

# interpolation

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# semester plan

Tu Nov 10 Least-squares and error

Th Nov 12 Case Study: Cancer Analysis

Tu Nov 17 **Building a basis for approximation (interpolation)**

Th Nov 19 non-linear Least-squares

Tu Dec 01 non-linear Least-squares

Th Dec 03 optimization methods

Tu Dec 08 Elements of Simulation + Review

# interpolation

Today's objectives:

1. Take a few points and **interpolate** instead of **fit**
2. Write the interpolant as a combination of \*basis\* functions
3. Implemente interpolation with several types of basis functions
4. Construct interpolation through a linear algebra problem

# interpolation: introduction

## Objective

Approximate an unknown function  $f(x)$  by an easier function  $g(x)$ , such as a polynomial.

## Objective (alt)

Approximate some data by a function  $g(x)$ .

Types of approximating functions:

1. Polynomials
2. Piecewise polynomials
3. Rational functions
4. Trig functions
5. Others (inverse, exponential, Bessel, etc)

# interpolation: introduction

How do we approximate  $f(x)$  by  $g(x)$ ? In what sense is the approximation a good one?

1. Least-squares:  $g(x)$  must deviate as little as possible from  $f(x)$  in the sense of a 2-norm: minimize  $\int_a^b |f(t) - g(t)|^2 dt$
2. Chebyshev:  $g(x)$  must deviate as little as possible from  $f(x)$  in the sense of the  $\infty$ -norm: minimize  $\max_{t \in [a,b]} |f(t) - g(t)|$ .
3. **Interpolation:  $g(x)$  must have the same values of  $f(x)$  at set of given points.**

# polynomial interpolation

Given  $n + 1$  distinct points  $x_0, \dots, x_n$ , and values  $y_0, \dots, y_n$ , find a polynomial  $p(x)$  of degree  $n$  so that

$$p(x_i) = y_i \quad i = 0, \dots, n$$

- A polynomial of degree  $n$  has  $n + 1$  degrees-of-freedom:

$$p(x) = a_0 + a_1x + \dots + a_nx^n$$

- $n + 1$  constraints determine the polynomial uniquely:

## Theorem

If points  $x_0, \dots, x_n$  are distinct, then for arbitrary  $y_0, \dots, y_n$ , there is a *unique* polynomial  $p(x)$  of degree at most  $n$  such that  $p(x_i) = y_i$  for  $i = 0, \dots, n$ .

# monomials

First attempt: try picking

$$p(x) = a_0 + a_1x + a_2x^2 + \cdots + a_nx^n$$

So for each  $x_i$  we have

$$p(x_i) = a_0 + a_1x_i + a_2x_i^2 + \cdots + a_nx_i^n = y_i$$

OR

$$a_0 + a_1x_0 + a_2x_0^2 + \cdots + a_nx_0^n = y_0$$

$$a_0 + a_1x_1 + a_2x_1^2 + \cdots + a_nx_1^n = y_1$$

$$a_0 + a_1x_2 + a_2x_2^2 + \cdots + a_nx_2^n = y_2$$

$$a_0 + a_1x_3 + a_2x_3^2 + \cdots + a_nx_3^n = y_3$$

$\vdots$

$$a_0 + a_1x_n + a_2x_n^2 + \cdots + a_nx_n^n = y_n$$

# monomial: the problem

$$\begin{bmatrix} 1 & x_0 & x_0^2 & \dots & x_0^n \\ 1 & x_1 & x_1^2 & \dots & x_1^n \\ 1 & x_2 & x_2^2 & \dots & x_2^n \\ & & & \vdots & \\ 1 & x_n & x_n^2 & \dots & x_n^n \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ a_2 \\ \vdots \\ a_n \end{bmatrix} = \begin{bmatrix} y_0 \\ y_1 \\ y_2 \\ \vdots \\ y_n \end{bmatrix}$$

## Question

- Is this a “good” system to solve?



# example

Consider Gas prices (in cents) for the following years:

x	year	1986	1988	1990	1992	1994	1996
y	price	133.5	132.2	138.7	141.5	137.6	144.2

```
1 year = np.array([1986, 1988, 1990, 1992, 1994, 1996])
2 price= np.array([133.5, 132.2, 138.7, 141.5, 137.6,
3                 144.2])
4 M = np.vander(year)
5 a = np.linalg.solve(M,price)
6
7 x = np.linspace(1986,1996,200)
8 p = np.polyval(a,x)
9 plt.plot(year,price,'o',x,p,'-')
```

# back to the basics...

## Example

Find the interpolating polynomial of least degree that interpolates

$x$	$1.4$	$1.25$
$y$	$3.7$	$3.9$

Directly

$$\begin{aligned} p_1(x) &= \left( \frac{x - 1.25}{1.4 - 1.25} \right) 3.7 + \left( \frac{x - 1.4}{1.25 - 1.4} \right) 3.9 \\ &= 3.7 + \left( \frac{3.9 - 3.7}{1.25 - 1.4} \right) (x - 1.4) \\ &= 3.7 - \frac{4}{3}(x - 1.4) \end{aligned}$$

What have we done? We've written  $p(x)$  as

$$p(x) = \left( \frac{x - x_1}{x_0 - x_1} \right) y_0 + \left( \frac{x - x_0}{x_1 - x_0} \right) y_1$$

- the sum of two linear polynomials
- the first is zero at  $x_1$  and 1 at  $x_0$
- the second is zero at  $x_0$  and 1 at  $x_1$
- these are the two linear Lagrange basis functions:

$$\ell_0(x) = \frac{x - x_1}{x_0 - x_1} \quad \ell_1(x) = \frac{x - x_0}{x_1 - x_0}$$

# lagrange

## Example

Write the Lagrange basis functions for

$$\begin{array}{c|ccc} x & \frac{1}{3} & \frac{1}{4} & 1 \\ \hline y & 2 & -1 & 7 \end{array}$$

Directly

$$\ell_0(x) = \frac{(x - \frac{1}{4})(x - 1)}{(\frac{1}{3} - \frac{1}{4})(\frac{1}{3} - 1)}$$

$$\ell_1(x) = \frac{(x - \frac{1}{3})(x - 1)}{(\frac{1}{4} - \frac{1}{3})(\frac{1}{4} - 1)}$$

$$\ell_2(x) = \frac{(x - \frac{1}{3})(x - \frac{1}{4})}{(1 - \frac{1}{3})(1 - \frac{1}{4})}$$

The general Lagrange form is

$$l_k(x) = \prod_{i=0, i \neq k}^n \frac{x - x_i}{x_k - x_i}$$

The resulting interpolating polynomial is

$$p(x) = \sum_{k=0}^n l_k(x) y_k$$

## example

Find the equation of the parabola passing through the points (1,6), (-1,0), and (2,12)

$$x_0 = 1, x_1 = -1, x_2 = 2; \quad y_0 = 6, y_1 = 0, y_2 = 12;$$

$$\begin{aligned}l_0(x) &= \frac{(x-x_1)(x-x_2)}{(x_0-x_1)(x_0-x_2)} = \frac{(x+1)(x-2)}{(2)(-1)} \\l_1(x) &= \frac{(x-x_0)(x-x_2)}{(x_1-x_0)(x_1-x_2)} = \frac{(x-1)(x-2)}{(-2)(-3)} \\l_2(x) &= \frac{(x-x_0)(x-x_1)}{(x_2-x_0)(x_2-x_1)} = \frac{(x-1)(x+1)}{(1)(3)}\end{aligned}$$

$$\begin{aligned}p_2(x) &= y_0 l_0(x) + y_1 l_1(x) + y_2 l_2(x) \\&= -3 \times (x+1)(x-2) + 0 \times \frac{1}{6}(x-1)(x-2) \\&\quad + 4 \times (x-1)(x+1) \\&= (x+1)[4(x-1) - 3(x-2)] \\&= (x+1)(x+2)\end{aligned}$$

## summary so far:

- Monomials:  $p(x) = a_0 + a_1x + \dots + a_nx^n$  results in poor conditioning
- Monomials: but evaluating the Monomial interpolant is cheap (nested iteration)
- Lagrange:  $p(x) = \ell_0(x)y_0 + \dots + \ell_n(x)y_n$  is very well behaved.
- Lagrange: but evaluating the Lagrange interpolant is expensive (each basis function is of the same order and the interpolant is not easily reduced to nested form)

## fixing monomials, fixing lagrange

Back to the gas price example. Suppose we use a better basis like

$$(x - \bar{x})^k$$

instead of

$$x^k$$

For example,  $\bar{x} = \text{average}(x_i), i = 0, \dots, n$ .

The basis  $(x - \bar{x})^k$  are called *shifted monomials* because  $x$  is shifted by  $\bar{x}$ .



# recall: monomials

Obvious attempt: try picking

$$p(x) = a_0 + a_1x + a_2x^2 + \cdots + a_nx^n$$

So for each  $x_i$  we have

$$p(x_i) = a_0 + a_1x_i + a_2x_i^2 + \cdots + a_nx_i^n = y_i$$

OR

$$\begin{bmatrix} 1 & x_0 & x_0^2 & \cdots & x_0^n \\ 1 & x_1 & x_1^2 & \cdots & x_1^n \\ & & & \vdots & \\ 1 & x_n & x_n^2 & \cdots & x_n^n \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ \vdots \\ a_n \end{bmatrix} = \begin{bmatrix} y_0 \\ y_1 \\ \vdots \\ y_n \end{bmatrix}$$

That is,

$$a = M^{-1}y$$

Very bad matrix: terribly ill-conditioned, inverse entries are **large**

Very bad evaluation: values are **huge**

# recall: lagrange

The general Lagrange form is

$$l_k(x) = \prod_{i=0, i \neq k}^n \frac{x - x_i}{x_k - x_i}$$

The resulting interpolating polynomial is

$$p(x) = \sum_{k=0}^n l_k(x) y_k$$

## example

Find the equation of a quadratic passing through the points (0,-1), (1,-1), and (2,7).

$$x_0 = 0, x_1 = 1, x_2 = 2 \quad y_0 = -1, y_1 = -1, y_2 = 7$$

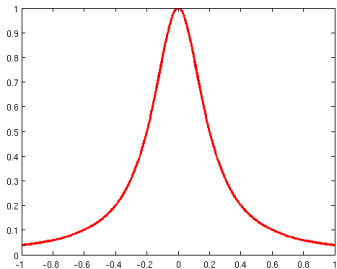
1. Form the Lagrange basis functions,  $l_i(x)$  with  $l_i(x_j) = \delta_{ij}$
2. Combine the Lagrange basis functions

$$\begin{aligned} p_2(x) &= y_0 l_0(x) + y_1 l_1(x) + y_2 l_2(x) \\ &= (-1) \frac{(x-1)(x-2)}{2} + (-1) \frac{x(x-2)}{-1} + (7) \frac{x(x-1)}{2} \end{aligned}$$

Evaluate is *nice*, but *expensive*: no easy nested form.

# how bad is polynomial interpolation?

Let's take something very smooth function



How does interpolation behave?

# some analysis...

what can we say about

$$e(t) = f(t) - p_n(t)$$

at some point  $x$ ? Consider  $p = 1$ : linear interpolation of a function at  $x = x_0, x_1$

- want: error at  $x$ ,  $e(x)$
- look at

$$g(t) = e(t) - \frac{(t-x_0)(t-x_1)}{(x-x_0)(x-x_1)}e(x)$$

- $g(t)$  is 0 at  $t = x_0, x_1, x$
- so  $g'(t)$  is zero at two points,  $g''(t)$  is zero at one point, call it  $c$

$$\begin{aligned}0 &= g''(c) = e''(t) - 2 \frac{e(x)}{(x-x_0)(x-x_1)} \\ &= f''(t) - 2 \frac{e(x)}{(x-x_0)(x-x_1)} \\ e(x) &= \frac{(x-x_0)(x-x_1)}{2} f''(c)\end{aligned}$$

## Theorem: Interpolation Error I

If  $p_n(x)$  is the (at most)  $n$  degree polynomial interpolating  $f(x)$  at  $n + 1$  distinct points and if  $f^{(n+1)}$  is continuous, then

$$e(x) = f(x) - p_n(x) = \frac{1}{(n+1)!} f^{(n+1)}(c) \prod_{i=0}^n (x - x_i)$$

## Theorem: Bounding Lemma

Suppose  $x_i$  are equispaced in  $[a, b]$  for  $i = 0, \dots, n$ . Then

$$\prod_{i=0}^n |x - x_i| \leq \frac{h^{n+1}}{4} n!$$

## Theorem: Interpolation Error II

Let  $|f^{(n+1)}(x)| \leq M$ , then with the above,

$$|f(x) - p_n(x)| \leq \frac{Mh^{n+1}}{4(n+1)}$$

We have two options:

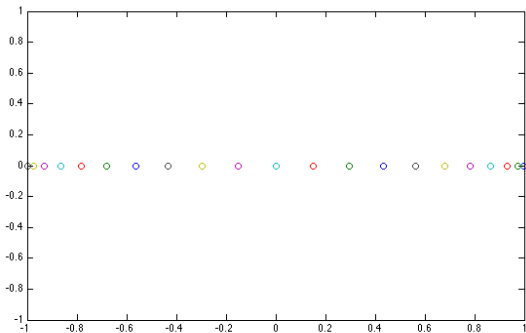
1. move the nodes: Chebychev nodes
2. piecewise polynomials (splines)

Option #1: Chebychev nodes in  $[-1, 1]$

$$x_i = \cos\left(\pi \frac{2i+1}{2n+2}\right), \quad i = 0, \dots, n$$

Option #2: piecewise polynomials...

# chebychev nodes

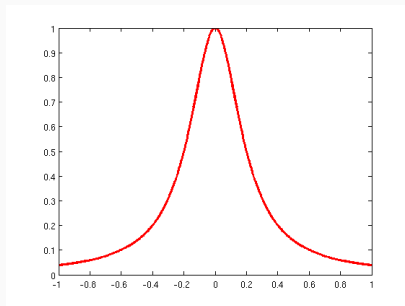


- Can obtain nodes from equidistant points on a circle projected down
- Nodes are non uniform and non nested



# chebychev nodes

High degree polynomials using equispaced points suffer from many oscillations



- Points are bunched at the ends of the interval
- Error is distributed more evenly