

CS 450: Numerical Analysis¹

Eigenvalue Problems

University of Illinois at Urbana-Champaign

¹*These slides have been drafted by Edgar Solomonik as lecture templates and supplementary material for the book “Scientific Computing: An Introductory Survey” by Michael T. Heath ([slides](#)).*

Eigenvalues and Eigenvectors

- ▶ A matrix A has eigenvector-eigenvalue pair (eigenpair) (λ, x) if

right-eigenpair if

$$Ax = \lambda x$$

if symmetric
=

left-eigenpair

$$x^T A = \lambda x^T$$

$$A^T x = \lambda x$$

eigenvector x

αx is also

$$\alpha Ax = \alpha \lambda x$$

1D eigenspace $\in \mathbb{C}^n$

$Ay = \lambda y$ then $A(x+y) = \lambda(x+y)$

- ▶ Each $n \times n$ matrix has up to n eigenvalues, which are either real or complex

if $a+bi \rightarrow$ eigenvalue
also $a-bi \rightarrow$ eigenvalue

up to n diff. eigenvalues

sum of multiplicities is $\leq n$

Eigenvalue Decomposition

- ▶ If a matrix A is diagonalizable, it has an *eigenvalue decomposition*

$$A = XDX^{-1}$$

right x -eigenvectors \swarrow X^{-1} is left eigenvectors (rows)

$$AX = XD = \begin{bmatrix} d_{11}x_1 & & \\ & d_{22}x_2 & \\ & & \dots & \\ & & & d_{nn}x_n \end{bmatrix}$$

if not diagonalizable, it is defective

- ▶ A and B are *similar*, if there exist Z such that $A = ZBZ^{-1}$

$$B = XDX^{-1}$$

similarity transformation

$$A = \underbrace{ZXDX^{-1}Z^{-1}}_{\text{eigenvectors}}$$

A and B are orthogonally similar if $\exists X$, s.t. $X^{-1} = X^T$, and $XAX^T = B$

If $A \in \mathbb{R}^{n \times n}$ normal, so $A^H A = A A^H$ } symmetric
if real $A^T A = A A^T$

$$\text{then } A = XDX^T = XDX^{-1}$$

orthogonally similar

$$A = U \Sigma V^T$$

$$A A^T = U \Sigma^2 U^T$$

$$A^T A = V \Sigma^2 V^T$$

if A is normal, then $U = V$, so $A = U \Sigma U^T$

} eigenvalue decompositions

property
symmetric
normal (real)
Hermitian
 $A = A^H$

similarity
orthogonally
orthogonally
unitarily
 $X = X^H$

matrix
real diagonal
real diagonal
real diagonal

Rayleigh Quotient and Power Iteration

- ▶ For any vector x , the *Rayleigh quotient* provides an estimate (lower-bound) on some eigenvalue λ of A :

$$\rho_A(x) = \frac{x^T A x}{x^T x} \quad \text{if } x \text{ is eigenvector } \rho_A(x) = \lambda$$

is not eigenvector $\rho_A(x) = \underset{\alpha}{\operatorname{argmin}} \|x\alpha - y\|_2$

normal eqns are $x^T x \alpha = x^T y$ Ax

$$\alpha = \frac{x^T y}{x^T x} = \frac{x^T A x}{x^T x}$$

- ▶ *Power iteration* can be used to compute the largest eigenvalue of a real symmetric matrix A :

pick a random $x^{(0)}$ and iterate $\left| \begin{array}{l} \text{linear convergence} \\ \text{rate } \lambda_1 / \lambda_2 \end{array} \right.$

$$x^{(i)} = A x^{(i-1)} = A(\alpha_1 y_1 + \alpha_2 y_2 + \dots) = \underbrace{\alpha_1 \lambda_1}_{\text{dominant}} y_1 + \alpha_2 \lambda_2 y_2 + \dots$$

if A is symmetric, eigvals $\lambda_1 > \lambda_2 > \dots$
eigvectors y_1, y_2, \dots

Inverse Iteration and Rayleigh Quotient Iteration

converge to eigenvalue closest to σ

- ▶ **Inverse iteration** uses LU/QR/SVD of A to run power iteration on A^{-1}

$$A^{-1} = X D^{-1} X^{-1}$$

$$\lambda_{\max}(A^{-1}) = 1/\lambda_{\min}(A)$$

shifted inverse iteration
decomposes $A - \sigma I$

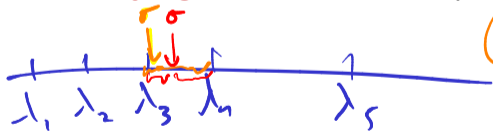
$$A = X D X^{-1}$$

$$A x^{(i+1)} = x^{(i)}$$

solve linear system at each step

given any factorization of A
solve costs $O(n^2)$

- ▶ **Rayleigh Quotient iteration** provides rapid convergence to an eigenpair



$$(A - \rho_A(x_i) I) x_{i+1} = x_i$$

solve

iterations
cost $O(n^3)$,
cubic convergence

$$\lambda_{\max}((A - \sigma I)^{-1}) = \frac{-1}{\lambda_3 - \sigma}$$

$$\lambda_{\text{2nd-max}}((A - \sigma I)^{-1}) = \frac{-1}{\lambda_4 - \sigma}$$

conv rate

$$\frac{\lambda_{\max}(\dots)}{\lambda_{\text{2nd-max}}(\dots)} = \frac{|\lambda_4 - \sigma|}{|\lambda_3 - \sigma|}$$

small

Perturbation Analysis of Eigenvalue Problems

- Suppose we seek eigenvalues $D = X^{-1}AX$, but find those of a slightly perturbed matrix $D + \delta D = \hat{X}^{-1}(A + \delta A)\hat{X}$: $\hat{X} = X + \delta X$

$$B = \hat{X}^{-1}A\hat{X}$$

$$B = \tilde{D} + \delta B$$

diagonal

$$\delta D = X^{-1}\delta S S X$$

$$\|\delta D\| \leq \|X^{-1}\| \|\delta S\| \|X\| = \kappa(X) \|\delta S\|$$

general perturbation, moreover $\text{eig}(B) = D$

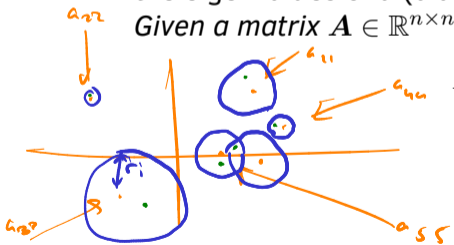
- Gershgorin's theorem allows us to bound the effect of the perturbation on the eigenvalues of a (diagonal) matrix:

Given a matrix $A \in \mathbb{R}^{n \times n}$, let $r_i = \sum_{j \neq i} |a_{ij}|$, define the Gershgorin disks as

$$D_i = \{z \in \mathbb{C} : |z - a_{ii}| \leq r_i\}.$$

sum of absolute values of i th row of A

eigenvalues diagonal entries



Conditioning of Particular Eigenpairs

- ▶ Consider the effect of a matrix perturbation on an eigenvalue λ associated with a right eigenvector x and a left eigenvector y^H , $\lambda = y^H A x / y^H x$

- ▶ A more accurate eigenvalue approximation than Rayleigh quotient for a normalized perturbed eigenvector (e.g. iterative guess) $\hat{x} = x + \delta x$, can be obtained with an estimate of both eigenvectors (also $\hat{y} = y + \delta y$),

Computing Eigenvectors of Matrices in Schur Form

- ▶ Given the eigenvectors of one matrix, we seek those of a similar matrix:

- ▶ Its easy to obtain eigenvectors of triangular matrix T :

Matrix Reductions

- ▶ Any matrix is orthogonally similar to an *upper-Hessenberg* (upper-triangular plus the first subdiagonal) matrix H , i.e. $A = QHQ^T$:

- ▶ In the symmetric case, Hessenberg form implies tridiagonal:

QR Iteration

- ▶ In orthogonal iteration $\hat{Q}_{i+1}\hat{R}_{i+1} = A\hat{Q}_i$, QR iteration computes $A_{i+1} = R_i Q_i = \hat{Q}_{i+1}^T A \hat{Q}_{i+1}$ at iteration i :

QR Iteration Complexity

- ▶ QR iteration is accelerated by first reducing to upper-Hessenberg or tridiagonal form:

Solving Tridiagonal Symmetric Eigenproblems

A rich variety of methods exists for the tridiagonal eigenproblem:

- ▶ QR iteration

- ▶ Divide and conquer

Solving the Secular Equation

To solve the eigenproblem at each step, the divide and conquer method needs to diagonalize a rank-1 perturbation of a diagonal matrix

$$\mathbf{A} = \mathbf{D} + \alpha \mathbf{u}\mathbf{u}^T$$

Introduction to Krylov Subspace Methods

- ▶ Define k -dimensional Krylov subspace matrix

$$\mathbf{K}_k = [\mathbf{x}_0 \quad \mathbf{A}\mathbf{x}_0 \quad \cdots \quad \mathbf{A}^{k-1}\mathbf{x}_0]$$

- ▶ Show that $\mathbf{K}_n^{-1}\mathbf{A}\mathbf{K}_n$ is a companion matrix \mathbf{C} :

Krylov Subspaces

- ▶ Given $QR = K_k$, we obtain an orthonormal basis for the Krylov subspace,

$$\mathcal{K}_k(\mathbf{A}, \mathbf{x}_0) = \text{span}(\mathbf{Q}) = \{\rho_{\mathbf{A}}(\mathbf{A})\mathbf{x}_0 : \text{deg}(\rho_{\mathbf{A}}) < k\}$$

- ▶ Consider whether $k - 1$ steps of power iteration starting from \mathbf{x}_0 lead to an approximation in the Krylov subspace, also consider QR (subspace) iteration:

Krylov Subspace Methods

- ▶ Given $QR = K_k$, we obtain an orthonormal basis for the Krylov subspace and $H_k = Q^T A Q$ which minimizes $\|AQ - QH\|_2$:

- ▶ H_k is Hessenberg, because the companion matrix C_k is Hessenberg:

Rayleigh-Ritz Procedure

- ▶ The eigenvalues/eigenvectors of \mathbf{H}_k are the *Ritz values/vectors*:
- ▶ The Ritz vectors and values are the *ideal approximations* of the actual eigenvalues and eigenvectors based on only \mathbf{H}_k and \mathbf{Q} :

Arnoldi Iteration

- ▶ Arnoldi iteration computes H directly using the recurrence $\mathbf{q}_i^T \mathbf{A} \mathbf{q}_j = h_{ij}$:

- ▶ After each matrix-vector product, orthogonalization is done with respect to each previous vector:

Convergence of Lanczos Iteration

- ▶ Cauchy interlacing theorem: eigenvalues of \mathbf{H}_k , $\tilde{\lambda}_1 \geq \dots \geq \tilde{\lambda}_n$ with respect to eigenvalues of \mathbf{A} , $\lambda_1 \geq \dots \geq \lambda_n$ satisfy

$$\lambda_i \leq \tilde{\lambda}_i \leq \lambda_{n-k+i}$$

- ▶ Convergence to extremal eigenvalues is generally fastest

Applications of Eigenvalue Problems: Matrix Functions

▶ Given $\mathbf{A} = \mathbf{X}\mathbf{D}\mathbf{X}^{-1}$ how can we compute \mathbf{A}^k ?

▶ What about $e^{\mathbf{A}}$? $\log(\mathbf{A})$? generally $f(\mathbf{A})$?

Applications of Eigenvalue Problems: Differential Equations

- ▶ Consider solutions to an ordinary differential equation of the form $\frac{d\mathbf{x}}{dt}(t) = \mathbf{A}\mathbf{x}(t) + \mathbf{f}(t)$ with $\mathbf{x}(0) = \mathbf{x}_0$:

$$\mathbf{x}(t) = e^{t\mathbf{A}}\mathbf{x}_0 + \int_0^t e^{(t-\tau)\mathbf{A}}\mathbf{f}(\tau)d\tau$$

- ▶ Using $\mathbf{A} = \mathbf{X}\mathbf{D}\mathbf{X}^{-1}$ permits us to compute the solution explicitly (Jordan form also suffices if \mathbf{A} is defective):

Differential Equations using the Generalized Eigenvalue Problem

- ▶ Consider a more general linear differential equation of the form $B \frac{dx}{dt}(t) = Ax(t) + f(t)$ with $x(0) = x_0$, which we can reduce to the usual form by premultiplying with B^{-1} :

- ▶ If we can find X such that $A = XD_A X^{-1}$ and $B = XD_B X^{-1}$ we could solve this equation while preserving symmetry of A and B :

Canonical Forms Generalized Eigenvalue Problem

- ▶ For nonsingular U, V , $A - \lambda B = U(J - \lambda I)V^T$ where J is in Jordan form

- ▶ For some unitary P, Q , $A = PT_AQ^H$ and $B = PT_BQ^H$ where T_A and T_B are triangular

Nonlinear Eigenvalue Problem

- ▶ In a polynomial eigenvalue problem, we seek solutions λ, \mathbf{x} to

$$\sum_{i=0}^d \lambda^i \mathbf{A}_i \mathbf{x} = \mathbf{0}$$

- ▶ Assuming for simplicity that $\mathbf{A}_d = \mathbf{I}$, solutions are given by solving the matrix eigenvalue problem with the block-companion matrix

$$\begin{bmatrix} -\mathbf{A}_{d-1} & \cdots & -\mathbf{A}_0 \\ \mathbf{I} & \mathbf{0} & \cdots \\ & \ddots & \ddots \end{bmatrix}$$