

CS 450: Numerical Analysis¹

Linear Least Squares

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¹*These slides have been drafted by Edgar Solomonik as lecture templates and supplementary material for the book “Scientific Computing: An Introductory Survey” by Michael T. Heath ([slides](#)).*

Linear Least Squares

- ▶ Find $\mathbf{x}^* = \operatorname{argmin}_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2$ where $\mathbf{A} \in \mathbb{R}^{m \times n}$.

- ▶ Given the SVD $\mathbf{A} = \mathbf{U}\mathbf{\Sigma}\mathbf{V}^T$ we have $\mathbf{x}^* = \mathbf{V}\mathbf{\Sigma}^\dagger\mathbf{U}^T\mathbf{b}$, where $\mathbf{\Sigma}^\dagger$ contains the reciprocal of all nonzeros in $\mathbf{\Sigma}$:

Normal Equations

- ▶ *Normal equations* are given by solving $A^T A x = A^T b$:

- ▶ However, solving the normal equations is a more ill-conditioned problem than the original least squares algorithm

Solving the Normal Equations

- ▶ If \mathbf{A} is full-rank, then $\mathbf{A}^T \mathbf{A}$ is symmetric positive definite (SPD):

- ▶ Since $\mathbf{A}^T \mathbf{A}$ is SPD we can use Cholesky factorization, to factorize it and solve linear systems:

QR Factorization

- ▶ If A is full-rank there exists an orthogonal matrix Q and a unique upper-triangular matrix R with a positive diagonal such that $A = QR$

- ▶ A reduced QR factorization (unique part of general QR) is defined so that $Q \in \mathbb{R}^{m \times n}$ has orthonormal columns and R is square and upper-triangular

- ▶ We can solve the normal equations (and consequently the linear least squares problem) via reduced QR as follows

Review: QR factorization

$$A \in \mathbb{R}^{m \times n}$$

$$\begin{bmatrix} \end{bmatrix} = \begin{bmatrix} \end{bmatrix} \begin{bmatrix} \end{bmatrix} = \begin{bmatrix} \end{bmatrix} \begin{bmatrix} \end{bmatrix}$$

\uparrow
orthogonal

$$\begin{bmatrix} \end{bmatrix} \begin{bmatrix} \end{bmatrix} = \begin{bmatrix} \end{bmatrix} \begin{bmatrix} \end{bmatrix}$$

reduced QR

$$Ax = b$$

$$QRx = b$$

$$Rx = Q^T b$$

$$\begin{bmatrix} \end{bmatrix} \begin{bmatrix} \end{bmatrix}$$

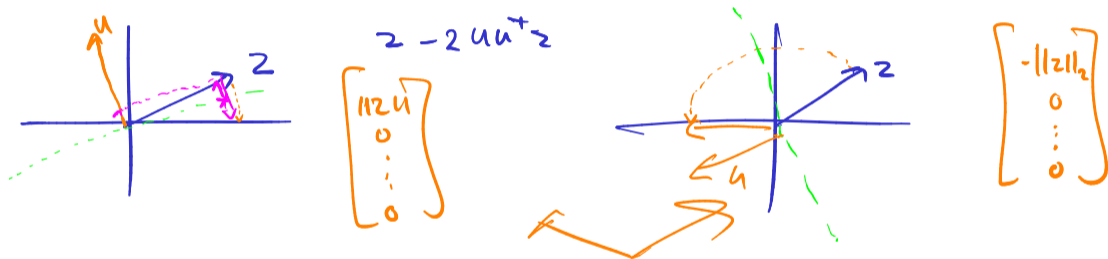
$$A^T Ax = A^T b$$

$$\cancel{R^T} \underbrace{Q^T Q}_I R x = \cancel{R^T} Q^T b$$

$$\Leftrightarrow Rx = Q^T b$$

Householder QR Factorization

- ▶ A Householder transformation $Q = I - 2uu^T$ is an orthogonal matrix defined to annihilate entries of a given vector z , so $\|z\|_2 Q e_1 = z$:



- ▶ Imposing this form on Q leaves exactly two choices for u given z ,

$$u = \frac{z \pm \|z\|_2 e_1}{\|z \pm \|z\|_2 e_1\|_2}$$

Applying Householder Transformations

- ▶ The product $x = Qw$ can be computed using $O(n)$ operations if Q is a Householder transformation

- ▶ Householder transformations are also called *reflectors* because their application reflects a vector along a hyperplane (changes sign of component of w that is parallel to u)

$$Q^T = Q$$

$$Q_1 A \Rightarrow \begin{bmatrix} \|a_1\| & u \\ 0 & B \end{bmatrix}$$

$$Q^T A = \text{closer to } R$$

$$\hat{Q}_2 B = \begin{bmatrix} \|b_1\| & v \\ 0 & \dots \end{bmatrix}$$

continue

$$A = QR$$

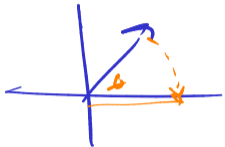
Householder reflector

$$Q^T = Q_1 \cdot Q_2 \cdot \dots \cdot Q_n \begin{bmatrix} I \\ \hat{a}_2 \\ \hat{a}_3 \end{bmatrix}$$

$$\begin{matrix} \begin{matrix} \left(I - 2uu^T \right) \\ \left(I - 2uu^T \right) v = v - 2u \langle u, v \rangle \end{matrix} \\ \begin{matrix} \downarrow n \times n \\ \Rightarrow \begin{matrix} O(n^2) \\ O(n) \end{matrix} \end{matrix} \end{matrix}$$

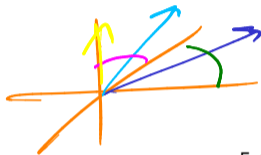
Givens Rotations

- ▶ Householder reflectors reflect vectors, Givens rotations rotate them



orthogonal matrices are compositions of

- rotations
- reflections



- ▶ Givens rotations are defined by orthogonal matrices of the form

$$\begin{bmatrix} c & s \\ -s & c \end{bmatrix}$$

not symmetric

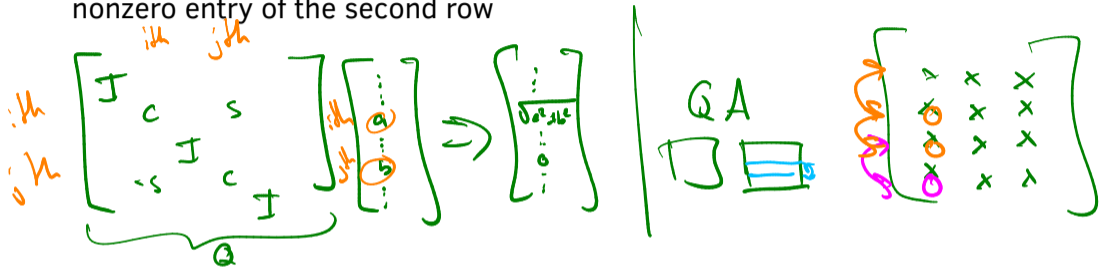
$$\begin{bmatrix} c & s \\ -s & c \end{bmatrix} \begin{bmatrix} a \\ b \end{bmatrix} = \begin{bmatrix} \sqrt{a^2 + b^2} \\ 0 \end{bmatrix}$$

$$c = \frac{a}{\sqrt{a^2 + b^2}}$$

$$s = \frac{b}{\sqrt{a^2 + b^2}}$$

QR via Givens Rotations

- ▶ We can apply a Givens rotation to a pair of matrix rows, to eliminate the first nonzero entry of the second row

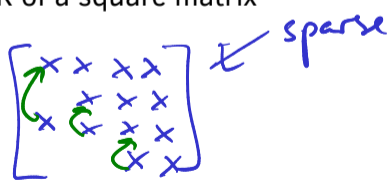


- ▶ Thus, $\frac{n(n-1)}{2}$ Givens rotations are needed for QR of a square matrix

$O(n)$ per Givens rotation

$O(n^3)$ in total

$\sim 2n^3$ in total vs $\frac{4}{3}n^3$ for Householder



Rank-Deficient Least Squares

- ▶ Suppose we want to solve a linear system or least squares problem with a (nearly) rank deficient matrix A

singular $\exists x \neq 0$ s.t. $Ax = 0$, $\sigma_{\min} = 0$
no unique soln to $Ax = b$ or $Ax \approx b$

if A is singular $f(A)$ probably isn't, but $\sigma_{\min} \approx \sigma_{\max} \epsilon_{\text{mach}}$

- ▶ Rank-deficient least squares problems seek a minimizer x of $\|Ax - b\|_2$ of minimal norm $\|x\|_2$ $\kappa(f(A)) = 1/\epsilon_{\text{mach}}$

$$x = A^+ b \Rightarrow \overline{A^+} = U \Sigma^+ U^T, \Sigma^+ = \begin{bmatrix} \sigma_1^+ & & \\ & \ddots & \\ & & \sigma_n^+ \end{bmatrix} \sigma_i^+ = \begin{cases} 1/\sigma_i & \text{if } \sigma_i > 0 \\ 0 & \text{if } \sigma_i = 0 \end{cases}$$

for numerical rank $\sigma_i^+ = 0$ if σ_i is very small

Truncated SVD

- ▶ After floating-point rounding, rank-deficient matrices typically regain full-rank but have nonzero singular values on the order of $\epsilon_{\text{mach}} \sigma_{\text{max}}$

alternative: regularization parameter α

(Tikhonov) $\min_x \|Ax - b\|_2 + \alpha \|x\|_2$

$$\begin{bmatrix} A \\ \alpha I \end{bmatrix} x \approx \begin{bmatrix} b \\ 0 \end{bmatrix} \Rightarrow \min_x \left\| \begin{bmatrix} Ax - b \\ \alpha x \end{bmatrix} \right\|_2$$

- ▶ By the Eckart-Young-Mirsky theorem, truncated SVD also provides the best low-rank approximation of a matrix (in 2-norm and Frobenius norm)

$A \approx \begin{bmatrix} \square \\ \square \end{bmatrix} \leftarrow$ low-rank approximation

rank- r truncated SVD is B , then $\|A - B\|_{2/F}$ is minimal

$$A = U \cdot \begin{bmatrix} \sigma_{\text{max}} & & \\ & \dots & \\ & & \sigma_6 & & \\ & & & & & 0 & & \end{bmatrix} U^T = \begin{bmatrix} \square \\ \square \end{bmatrix} \setminus \begin{bmatrix} \square \\ \square \end{bmatrix} = \begin{bmatrix} \square \\ \square \end{bmatrix} \setminus \begin{bmatrix} \square \\ \square \end{bmatrix}$$

QR with Column Pivoting

- QR with column pivoting provides a way to approximately solve rank-deficient least squares problems and compute the truncated SVD

$$\begin{array}{l}
 QR = AP \\
 \uparrow \\
 \text{permutation} \\
 \text{matrices}
 \end{array}
 \quad
 QR = \begin{array}{|l} \hline \text{ } \\ \hline \end{array}
 \left|
 \begin{array}{l}
 Ax \approx b \\
 QR P^T x \approx b \\
 \underbrace{y}_{Ry = Q^T b}
 \end{array}
 \right|
 \begin{array}{|l} \hline \text{ } \\ \hline \end{array}
 = AP$$

- A pivoted QR factorization can be used to compute a rank- r approximation

taking first r columns of AP , and compute SVD

for each column (Householder reflector)
 pivot column of largest norm to be leading column
 P_i $Q_i A P_i = \begin{bmatrix} \alpha & b \\ 0 & B \end{bmatrix} \rightarrow$ continue with B | Greedy algorithm

$Q(n \times r^2)$