

CS 450: Numerical Analysis¹

Numerical Integration and Differentiation

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¹*These slides have been drafted by Edgar Solomonik as lecture templates and supplementary material for the book “Scientific Computing: An Introductory Survey” by Michael T. Heath ([slides](#)).*

Integrability and Sensitivity

- ▶ Seek to compute $\mathcal{I}(f) = \int_a^b f(x)dx$:

- ▶ The condition number of integration is bounded by the distance $b - a$:

Quadrature Rules

- ▶ Approximate the integral $\mathcal{I}(f)$ by a weighted sum of function values:

$$\mathcal{I}(f) = \int_a^b f(x) dx \approx Q_n = \sum_{i=1}^n w_i \underbrace{f(x_i)}_{y_i} = \langle w, y \rangle$$

\uparrow
 n -point
 x_1, \dots, x_n
 nodes

- ▶ For a fixed set of n nodes, polynomial interpolation followed by integration give $(n-1)$ -degree quadrature rule:

$$\mathcal{I}(f) \approx Q_n = \mathcal{I}(p) = \mathcal{I}(\langle \{e_i\}_{i=1}^n, c \rangle) = \langle \underbrace{V^{-T} \{ \mathcal{I}(e_i) \}_{i=1}^n}_w, y \rangle$$

$V^{-1} y = \langle \{ \mathcal{I}(e_i) \}_{i=1}^n, V^{-1} y \rangle$
 \downarrow

degree $n-1$ polynomial interpolant of (x_1, y_1)
 \vdots
 (x_n, y_n)

$$p(x) = \sum_{i=1}^n c_i e_i(x) \Rightarrow \int_a^b p(x) dx = \sum_{i=1}^n c_i \int_a^b e_i(x) dx$$

Determining Weights in a General Basis

- ▶ A quadrature rule provides x and w so as to approximate

$$Q_n = \langle w, f \rangle$$

eqn.-spaced

Chebyshev points

find optimal distribution of points,
to maximize degree

- ▶ *Method of undetermined coefficients* obtains y from *moment equations* based on Vandermonde system:

weights w are independent of f or y
and so are valid provided interval of
integration is the same

Error in Newton-Cotes Quadrature

- Consider the Taylor expansion of f about the midpoint of the integration interval $m = (a+b)/2$:

$$f(x) = f(m) + f'(m)(x-m) + \frac{f''(m)}{2}(x-m)^2 + \dots$$

Integrating the Taylor approximation of f , we note that the odd terms drop:

$$I(f) = \underbrace{(b-a)f(m)}_{M(f)} + \underbrace{\frac{f''(m)}{24}(b-a)^3}_{E(f)} + O((b-a)^5)$$

midpoint rule is third order in accuracy
 → to bound error in Trapezoid rule
 want $f(m)$ in terms of $f(a)$, $f(b)$

Error Estimation

- ▶ The trapezoid rule is also first degree, despite using higher-degree polynomial interpolant approximation, since

$$f(m) = \frac{1}{2} \left[f(a) - \cancel{f'(m)(a-m)} - \frac{f''(m)}{2} (a-m)^2 - \dots \right. \\ \left. + f(b) - \cancel{f'(m)(b-m)} - \frac{f''(m)}{2} (b-m)^2 - \dots \right]$$

$\frac{b-a}{2}$
 $\frac{b-a}{2}$

- ▶ The above derivation allows us to obtain an error approximation via a difference of midpoint and trapezoidal rules:

$$I(f) = \underbrace{(b-a) \left[\frac{1}{2} (f(b) + f(a)) \right]}_{T(f)} - \frac{f''(m) (b-a)^2}{12} - O(b^2)$$

$$3E(f) = \underline{M(f)} - \underline{T(f)} + O((b-a)^2) \underbrace{2E(f)}$$

Error in Polynomial Quadrature Rules

- ▶ We can bound the error for ~~a~~ an arbitrary polynomial quadrature rule by

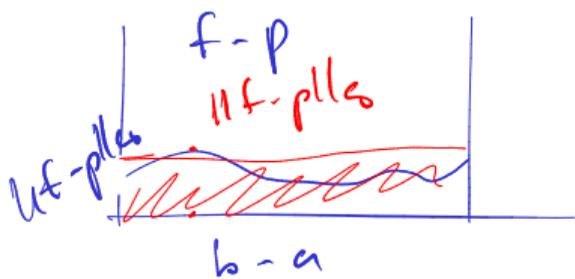
$$|I(f) - Q_n(f)| = |I(f) - I(p)|$$

$$= |I(f-p)|$$

$$\leq (b-a) \|f-p\|_\infty$$

$$\leq (b-a) \frac{(b-a)^n |f^{(n)}|_\infty}{n!}$$

$$= O((b-a)^{n+1})$$



polynomial interpolant

Conditioning of Newton-Cotes Quadrature

- ▶ We can ascertain stability of quadrature rules, by considering the amplification of a perturbation $\hat{f} = f + \delta f$:

$$Q_n(\hat{f}) - Q_n(f) = \langle w, \hat{y} \rangle - \langle w, y \rangle$$
$$= \langle w, \hat{y} - y \rangle$$

condition: $\sum_{i=1}^n w_i = b-a$

want to minimize $\|w\|, = \sum_{i=1}^n |w_i| \geq b-a$

$$= Q_n(\delta f) \leq \underbrace{\|w\|}_{\text{bad}}, \underbrace{\|\delta f\|_{\max}}$$

- ▶ Newton-Cotes quadrature rules have at least one negative weight for any $n \geq 11$:

So Newton-cotes quadrature rules can be ill-conditioned (suboptimal) although for $n=1,2,3$ they are fine

$\hat{y}_i = \hat{f}(x_i)$
 $y_i = f(x_i)$

Clenshaw-Curtis Quadrature

- ▶ To obtain a more stable quadrature rule, we need to ensure the integrated interpolant is well-behaved as n increases:

Chebyshev-spaced nodes

↳ interpolation error is smaller
↳ can show that $w_i > 0$

Gaussian Quadrature

- ▶ So far, we have only considered quadrature rules based on a fixed set of nodes, but we may also be able to choose nodes to maximize accuracy:

extra n parameters, for a total of $(n \text{ weights} \& \text{ nodes})$
 $2n$ degrees of freedom
 degree $2n-1$ quadrature rules
 picks x, w

- ▶ The **unique** n -point **Gaussian quadrature rule** is defined by the solution of the nonlinear form of the moment equations in terms of both x and w :

solve for x and w $V(x, \{e_j^n\}_{j=1}^n)^T w = \{I(e_j)\}_{j=1}^n$

independent of $y(f)$, dependent only on n, a, b

- solution generally exists and is unique
- $w_i > 0$, well-conditioned

Using Gaussian Quadrature Rules

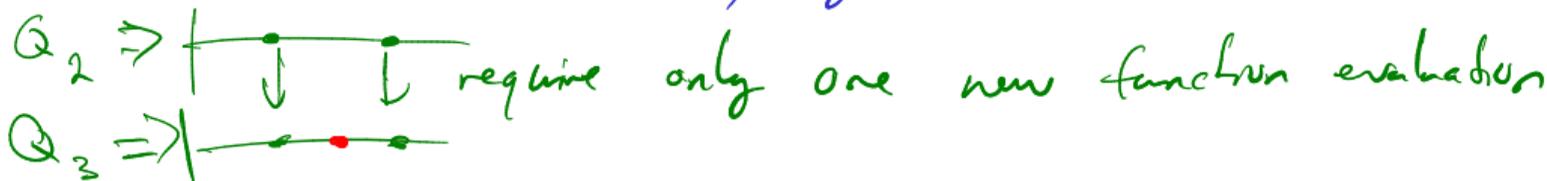
- ▶ Gaussian quadrature rules are hard to compute, but can be enumerated for a fixed interval, e.g. $a = 0, b = 1$, so it suffices to transform the integral to $[0, 1]$

$$I(f) = \int_a^b f(x) dx = \frac{1}{b-a} \int_0^1 g(t) dt \quad \begin{matrix} t = \frac{x-a}{b-a} \\ dt = \frac{1}{b-a} \end{matrix} \quad g(t) = f(t(b-a) + a)$$

can do similar to get to $[-1, 1]$

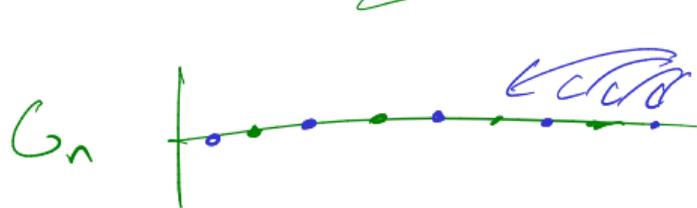
- ▶ Gaussian quadrature rules are accurate and stable but not progressive (nodes cannot be reused to obtain higher-degree approximation):

progressive rules reuse nodes, e.g.



Progressive Gaussian-like Quadrature Rules

- ▶ **Kronod** quadrature rules construct $(2n+1)$ -point $(3n+1)$ -degree quadrature K_{2n+1} that is progressive with respect to Gaussian quadrature rule G_n .



given G_n computing K_{2n+1} requires $n+1$ func. eval.
computing G_{2n+1} would require $2n+1$ degree $4n+1$

- ▶ **Patterson** quadrature rules use $2n+2$ more points to extend $(2n+1)$ -point Kronod rule to degree $6n+4$, while reusing all $2n+1$ points.
- ▶ Gaussian quadrature rules are in general open, but **Gauss-Radau** and **Gauss-Lobatto** rules permit including end-points:

Gauss rules are open
Gauss-Lobatto is closed

Gauss-Radau contains
1-endpoint

Composite and Adaptive Quadrature

- ▶ **Composite quadrature rules** are obtained by integrating a piecewise polynomial interpolant of f :

main advantage is accuracy \Rightarrow shrink $b-a$
(interval of integration)

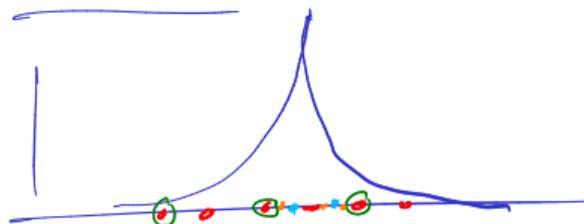
$$I(f) \approx I(p) = \sum_{i=1}^n \int_{x_i}^{x_{i+1}} p_i(x) dx$$

\uparrow
 $x_0 = a$
 $x_n = b$
 \uparrow
 i th piece

$$E(f) \approx O((b-a)^n)$$

\uparrow
error

- ▶ Composite quadrature can be done with adaptive refinement:

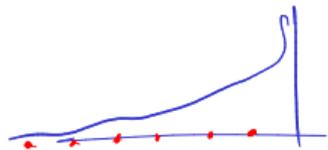


error near quickly changing part of function
converges slower (estimate says its large)

\rightarrow e.g. $E(f) \approx \frac{1}{3}(M(f) - T(f))$

More Complicated Integration Problems

- ▶ To handle improper integrals can either transform integral to get rid of infinite limit or use appropriate open quadrature rules.



- ▶ Double integrals can simply be computed by successive 1-D integration.

if can decouple

- ▶ High-dimensional integration is often effectively done by **Monte Carlo**:

$$I(f) \approx \frac{\Omega}{N} \sum_{i=1}^N f(x_i)$$

integration volume Ω N = number of samples N

x_i are sample from domain

error converges with $O(1/\sqrt{N})$

Integral Equations

- ▶ Rather than evaluating an integral, in solving an *integral equation* we seek to compute the integrand. A typical linear integral equation has the form

$$\int_a^b K(s, t)u(t)dt = f(s), \quad \text{where } K \text{ and } f \text{ are known.}$$

kernel known unknown function values known

linear system matrix

$$\int_a^b k(s, t) u(t) dt \approx \sum_{i=1}^n w_i k(s_i, t_j) u(t_j) dt = f(s_i)$$

- ▶ Using a quadrature rule with weights w_1, \dots, w_n and nodes t_1, \dots, t_n obtain

Extrapolation Techniques

- ▶ Given a series of approximate solutions produced by an iterative procedure, a more accurate approximation may be obtained by *extrapolating* this series.

- ▶ In particular, given two guesses, *Richardson extrapolation* eliminates the leading order error term:

