

- Exan let 1 :



Changing Condition Numbers

Once we have a matrix A in a linear system $A\mathbf{x} = \mathbf{b}$, are we stuck with its condition number? Or could we improve it?

$$\begin{array}{l} D A \mathbf{x} = D \mathbf{b} \\ M A \mathbf{x} = M \mathbf{b} \end{array} \quad \begin{array}{l} \text{non-sing. diag } D \\ \text{"left preconditioning"} \\ \hline \mathbf{x} = M \mathbf{y} \\ (A M) \mathbf{y} = \mathbf{b} \end{array}$$

right precond:

What is this called as a general concept?

preconditioning
common examples: rescale rows / cols

Not a matrix norm: Frobenius

The 2-norm is very costly to compute. Can we make something simpler?

$$\|A\|_F = \sqrt{\sum_{i=1}^m \sum_{j=1}^n |a_{ij}|^2}$$

What about its properties?

- definiteness
- semilinearity
- triangle ineq
- submultiplicative

Frobenius Norm: Properties

bit.ly/cs450-f22

Is the Frobenius norm induced by any vector norm?

$$\|I\|_F = \sqrt{n}$$

For any induced norm, $\|I\| = 1$

How does it relate to the SVD?

$$\|A\|_F = \|\sigma\|_2$$

$$\|A\| \|A^{-1}\| \geq \|AA^{-1}\| = \|I\| = 1$$

$$\begin{pmatrix} \frac{1}{2} & 0 \\ 0 & \frac{1}{4} \end{pmatrix} \quad \checkmark$$

$$\left(\begin{array}{c} \boxed{a_1} \\ \boxed{a_2} \end{array} \right) \left(\begin{array}{c} x_1 \\ x_2 \end{array} \right) = x_1 \vec{a}_1 + x_2 \vec{a}_2$$

$$|x_1| + |x_2| = 1$$

Solving Systems: Simple cases

Solve $D\mathbf{x} = \mathbf{b}$ if D is diagonal. (Computational cost?)

$$x_i = b_i / d_{ii} \quad O(n)$$

Solve $Q\mathbf{x} = \mathbf{b}$ if Q is orthogonal. (Computational cost?)

$$\vec{x} = Q^T \vec{b} \quad O(n^2)$$

Given SVD $A = U\Sigma V^T$, solve $A\mathbf{x} = \mathbf{b}$. (Computational cost?)

$$\begin{aligned} \text{Compute } z &= U^T b \\ \text{Solve } \Sigma y &= z \\ \text{Compute } \vec{x} &= V y \end{aligned} \quad \rightarrow \quad U \Sigma V^T x = b$$

Solving Systems: Triangular matrices

Solve

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ & a_{22} & a_{23} & a_{24} \\ & & a_{33} & a_{34} \\ & & & a_{44} \end{bmatrix} \begin{bmatrix} x \\ y \\ z \\ w \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ b_4 \end{bmatrix} \cdot \uparrow \text{ back subst.}$$

Fwd
subst



$$a_{44} w = b_4$$
$$a_{33} z + a_{34} w = b_3 \Leftrightarrow z = (b_3 - a_{34} w) / a_{33}$$

Cost: $O(n^2)$

Demo: Coding back-substitution [cleared]

What about non-triangular matrices?

Gaussian elim.

In-Class Activity: Matrix Norms and Conditioning

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Gaussian Elimination

Demo: Vanilla Gaussian Elimination [cleared]

What do we get by doing Gaussian Elimination?

$$A|b$$
$$\left(\begin{array}{c|c} & \\ & \\ & \\ & \end{array} \right)$$

Row Echelon Form

How is that different from being upper triangular?

$$\left(\begin{array}{c} \text{row} \\ \text{echelon} \\ \text{rank-revealing} \end{array} \right) \left(\begin{array}{c} \text{row} \\ \text{echelon} \end{array} \right)$$

What if we do not just eliminate downward but also upward?

Gauss-Jordan elim : more expensive than LU

LU Factorization

What is the LU factorization?

$$A = LU$$

Lower
Upper

$$\begin{matrix} \triangle L \\ \nabla U \end{matrix}$$

Solving $A\mathbf{x} = \mathbf{b}$

Does LU help solve $A\mathbf{x} = \mathbf{b}$?

$$A\mathbf{x} = \mathbf{b}$$

$$L\underbrace{U}_{\mathbf{y}}\mathbf{x} = \mathbf{b}$$

$$L\mathbf{y} = \mathbf{b}$$

$$U\mathbf{x} = \mathbf{y}$$

Determining an LU factorization

L of size $n \times n$ \rightarrow

$$\begin{pmatrix} a_{11} & a_{12}^T \\ \vec{a}_{21} & A_{22} \end{pmatrix} = \begin{pmatrix} L_{11} & \\ L_{21} & L_{22} \end{pmatrix} \begin{pmatrix} u_{11} & u_{12} \\ & U_{22} \end{pmatrix}$$

(Annotations: L_{11} is a scalar, L_{21} is a vector)

$$\begin{pmatrix} u_{11} & u_{12}^T \\ 0 & U_{22} \end{pmatrix} \Bigg| u$$

$$\begin{pmatrix} 1 & 0 \\ \vec{l}_{21} & L_{22} \end{pmatrix} \begin{pmatrix} a_{11} & a_{12}^T \\ \vec{a}_{21} & A_{22} \end{pmatrix}$$

- $a_{11} = u_{11}$
- $\vec{u}_n^T = \vec{a}_{12}^T$
- $\vec{a}_{21} = u_{11} \cdot \vec{l}_{21} \Leftrightarrow \vec{l}_{21} = \vec{a}_{21} / u_{11}$
- $A_{22} = \vec{l}_{21} u_{12}^T + L_{22} U_{22}$
- $L_{22} U_{22} = A_{22} - \vec{l}_{21} u_{12}^T$

\rightarrow LU of size $(n-1) \times (n-1)$

Demo: LU Factorization [cleared]

Computational Cost

What is the computational cost of multiplying two $n \times n$ matrices?

- ▶ $u_{11} = a_{11}, \mathbf{u}_{12}^T = \mathbf{a}_{12}^T.$
- ▶ $l_{21} = \mathbf{a}_{21}/u_{11}.$
- ▶ $L_{22}U_{22} = A_{22} - l_{21}\mathbf{u}_{12}^T.$

What is the computational cost of carrying out LU factorization on an $n \times n$ matrix?

[Demo: Complexity of Mat-Mat multiplication and LU \[cleared\]](#)