

QR Iteration/QR Algorithm

Orthogonal iteration:

$$X_0 = A$$

$$Q_k R_k = X_k$$

$$X_{k+1} = A Q_k$$

QR iteration:

$$\bar{X}_0 = A$$

$$\bar{Q}_k \bar{R}_k = \bar{X}_k \quad \left. \vphantom{\bar{Q}_k \bar{R}_k = \bar{X}_k} \right\} n \text{ times}$$

$$\bar{X}_{k+1} = \bar{R}_k \bar{Q}_k$$

Tracing through reveals:

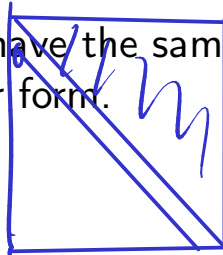
- $\hat{X}_k = \bar{X}_{k+1}$
- $Q_0 = \bar{Q}_0$
 $Q_1 = \bar{Q}_0 \bar{Q}_1$
 $Q_k = \bar{Q}_0 \bar{Q}_1 \cdots \bar{Q}_k$

- conv. slowly
 - expensive

Orthogonal iteration showed: $\hat{X}_k = \bar{X}_{k+1}$ converge. Also:

$$\bar{X}_{k+1} = \bar{R}_k \bar{Q}_k = \bar{Q}_k^T \bar{X}_k \bar{Q}_k,$$

so the \bar{X}_k are all similar \rightarrow all have the same eigenvalues.
 \rightarrow QR iteration produces Schur form.

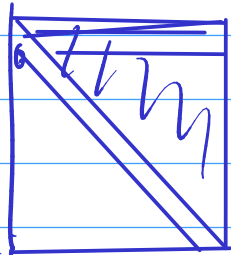
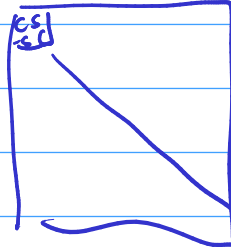
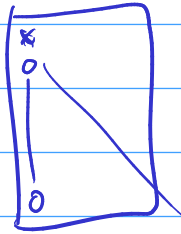


$$\left(\begin{array}{|c|} \hline | \\ \hline | \\ \hline | \\ \hline | \\ \hline | \\ \hline | \\ \hline | \\ \hline \end{array} \right) = QR$$

A

$$I - 2 \frac{uv^T}{v^T v}$$

$$H_n = H_1 A =$$



1 Householder: $O(n^2)$

n Householders $O(n^3)$

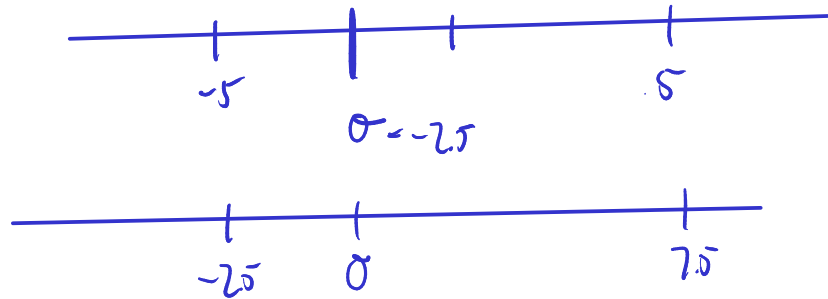
1 Givens rot $O(n)$

n Givens rot $O(n^2)$

QR Iteration: Incorporating a Shift

- How can we accelerate convergence of QR iteration using shifts?

$$\begin{aligned}\tilde{X}_0 &= A \\ \tilde{Q}_k \tilde{R}_k &= \tilde{X}_k - \sigma_k I_k \\ \tilde{X}_{k+1} &= \tilde{R}_k \tilde{Q}_k + \sigma_k I\end{aligned}$$



QR Iteration: Computational Expense

- A full QR factorization at each iteration costs $O(n^3)$ —can we make that cheaper?

Upper Hessenberg form

$$A = Q \begin{matrix} \square \\ \diagdown \\ \square \end{matrix} Q^T$$

By getting A into upper H. form before we start QR iteration, we can get the cost of 1 step of QR it. to $O(n^2)$.

$$A = Q \begin{matrix} \square \\ \diagdown \\ \square \end{matrix} Q^T$$

Demo: Householder Similarity Transforms

4.4 Krylov Space Methods

Q2 it.

$$\text{span} [A^l y_1, A^l y_2, \dots, A^l y_n]$$

Krylov space

$$[\underbrace{x}_x, \underbrace{Ax}_x, \underbrace{A^2x}_x, A^3x, \dots, A^{n-1}x]$$

$$K_k = \begin{pmatrix} | & & | \\ x_0 & \dots & x_{k-1} \\ | & & | \end{pmatrix} \quad (n \times k) \quad (1 \leq k \leq n)$$

$$AK_k = \begin{pmatrix} | & | & \dots & | & | \\ x_1 & x_2 & \dots & x_{k-1} & x_k \\ | & | & & | & | \end{pmatrix} = K_k \begin{pmatrix} 0 & 0 & & & \\ 1 & 0 & & & \\ 0 & 1 & & & \\ \vdots & \vdots & \ddots & \ddots & \\ 0 & & & 0 & K_{k-1}^{-1} x_k \end{pmatrix}$$

C_n

$$A K_n = K_n C_n$$

$$K_n^{-1} A K_n = C_n$$

Krylov space methods: Intro

- What subspaces can we use to look for eigenvectors?

Conditioning in Krylov Space Methods/Arnoldi Iteration

- What is a problem with Krylov space methods? How can we fix it?

$$Q_n R_n = K_n \quad \Rightarrow \quad Q_n = K_n R_n^{-1}$$

Then:

$$Q_n^T A Q_n = \underbrace{R_n^{-1} K_n^T A K_n R_n^{-1}}_{C_n}$$

$\underbrace{\hspace{10em}}_{C_n}$

$\underbrace{\hspace{10em}} = H$

$$Q_n^T A Q_n = H$$

$$A Q_n = Q_n H$$

$$Q_n = (q_1 \dots q_n)$$

$$A q_k = h_{1k} \vec{q}_1 + \dots + h_{k+1,k} \vec{q}_{k+1}$$

$$H = (h_{ij})$$

$$q_j^T A q_k = h_{jk}$$

$$q_i^T q_l = \begin{cases} 1 & i=l \\ 0 & i \neq l \end{cases}$$

Demo: Arnoldi Iteration (Part I)

Krylov: What about eigenvalues?

- How can we use Arnoldi/Lancos to compute eigenvalues?