

CS 450: Numerical Analysis

Lecture 7

Chapter 3 – Linear Least Squares

Stability and Efficiency of Linear Least Squares Algorithms

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Conditioning of Linear Least Squares

- ▶ Consider fitting a line to a collection of points, then perturbing the points:
 - ▶ *If our line closely fits all of the points, a small perturbation to the points will not change the ideal fit line (least squares solution) much. Note that, if a least squares solution has a very small residual, any other solution with a residual close to as small, should be close to parallel to this solution.*
 - ▶ *When the points are distributed erratically and do not admit a reasonable linear fit, then the least squares solution has a large residual, and totally different lines may exist with a residual nearly as small. For example, if the points are in a ball around the origin, any linear fit has the same residual. A tiny perturbation could then perturb the least squares solution to be perpendicular to the original.*
- ▶ LLS is ill-posed for any \mathbf{A} , unless we consider solving for a particular \mathbf{b}
 - ▶ *If \mathbf{b} is outside the span \mathbf{A} then any perturbation to \mathbf{A} or \mathbf{b} can completely defines the new solution. Similarly, if most of \mathbf{b} is outside the span of \mathbf{A} , a perturbation can cause the solution to fluctuate wildly.*
 - ▶ *On other hand, if for a particular \mathbf{b} we can find a solution with (near-)zero residual, a small relative perturbation to \mathbf{b} or \mathbf{A} will have an effect similar to that of a linear system perturbation (growth bounded by $\kappa(\mathbf{A}) = \sigma_{max}/\sigma_{min}$).*

Stability of Normal Equations

- ▶ Normal equations solve $\mathbf{A}^T \mathbf{A} \mathbf{x} = \mathbf{A} \mathbf{b}$:

Consider perturbation $\delta \mathbf{b}$ to \mathbf{b} , in the worst case $\mathbf{A} \mathbf{b}$ is perturbed by $\sigma_{\max} \delta \mathbf{b}$. Further $\kappa(\mathbf{A}^T \mathbf{A}) = \kappa(\mathbf{A})^2$ so the relative error can be amplified by the square of the expected condition number, given a small residual in the true solution.

- ▶ We can also obtain a QR factorization by means of Cholesky of $\mathbf{A}^T \mathbf{A}$ (known as Cholesky-QR):

- ▶ *The QR factorization $\mathbf{A} = \mathbf{Q} \mathbf{R}$ so long as \mathbf{R} has a positive diagonal, so we have*

$$\mathbf{A}^T \mathbf{A} = \mathbf{R}^T \mathbf{Q}^T \mathbf{Q} \mathbf{R} = \mathbf{R}^T \mathbf{R}$$

and thus can obtain \mathbf{R} by Cholesky $\mathbf{A}^T \mathbf{A}$, which is symmetric positive definite.

- ▶ *However, if condition number of $\mathbf{A}^T \mathbf{A}$ is high, we can lose positive-definiteness due to roundoff error and Cholesky factorization can break down.*
- ▶ *Moreover, if we obtain $\hat{\mathbf{Q}}$ via $\mathbf{A} \mathbf{R}^{-1}$, round-off error can result in a loss of orthogonality in the columns of $\hat{\mathbf{Q}}$.*
- ▶ *So long as $\kappa(\mathbf{A}) < \sqrt{1/\epsilon_{\text{mach}}}$, a second iteration of Cholesky-QR on the computed reduced $\hat{\mathbf{Q}} = \mathbf{Q} \hat{\mathbf{R}}$ obtains an orthogonal \mathbf{Q} and a correction $\hat{\mathbf{R}}$ to \mathbf{R} .*