CS 450: Numerical Anlaysis Lecture 10 Chapter 4 – Eigenvalue Problems Theory of Eigenvalue Solvers

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February 16, 2018

#### Perturbation Analysis of Eigenvalue Problems

Suppose we seek eigenvalues  $D = X^{-1}AX$ , but find those of a slightly perturbed matrix  $\boldsymbol{D} + \boldsymbol{\delta} \boldsymbol{D} = \boldsymbol{\hat{X}}^{-1} (\boldsymbol{A} + \boldsymbol{\delta} \boldsymbol{A}) \boldsymbol{\hat{X}}$ :  $D + \delta D = \hat{x}^{-1}A\hat{x} + \hat{x}^{-1}SA\hat{x}$  $= X^{-1}(A^{\perp}SA)X = X^{-1}AX + X^{\perp}SAX$  $eig(A + SA) = D + SD = eig(D + x - 1 SA \times)$ 11x - (A X) 5 R(X) | SAI = 11×-111 11×11 11 \$ A /1

(matrix is diagonal) D = A

L-X

for some perturbation (non-diagonal) SA  

$$eig(A : SA) = D + SD$$
 (A) (44er from D  
 $D + SD = \hat{X}^{-1}(A : SA \setminus \hat{X})$   
 $\hat{X} = J + SX \mid SD = \hat{X}^{-1}(SAX - SXA + ASX)$   
 $\hat{X} \approx J - SX \mid SD = \hat{X}^{-1}(SAX - SXA + ASX)$   
 $\hat{X} \approx J - SX \mid SD = \hat{X}^{-1}(SAX - SXA + ASX)$ 

### **Gershgorin Theorem**

A = D + O

Another way to show that the eigenvalues of a matrix are insensitive to perturbation is via Gershgorin theorem, which states that

 $\lambda(A) \in \bigcup_{A:st} p(A:s, r;)$ 

 $r_{i} = \hat{\mathcal{E}} |O_{ij}|$ 

G

And Ass

# **Conditioning of Particular Eigenpairs**

A=X"NX

 $\triangleright$  Consider the effect of a matrix perturbation on an eigenvalue  $\lambda$  associated with a right eigenvector x and a left eigenvector  $y^H$ ,  $\lambda = y^H A x / y^H x$ 

J <4, 1A + FA' Connect the notion of the angle between left and right eigenvectors to the magnitude of off-diagonal entries in the Schur form

<4, Ax > = <

y, x> < < y, 2A = < y, 2A

Orthogonal Iteration via QR Iteration  
In orthogonal iteration 
$$\hat{Q}_{i+1}\hat{R}_{i+1} = A\hat{Q}_i$$
 QR iteration computes  
 $A: G_1 P_1 A_{i+1} = R_i Q_i = \hat{Q}_{i+1}^T A \hat{Q}_{i+1}$  at iteration i:  
 $U = \lim_{i \to \infty} A_i = \lim_{i \to \infty} \hat{Q}_{i+1} A \hat{Q}_i$   
 $i \to \infty$   
 $A: = \lim_{i \to \infty} A_i = \lim_{i \to \infty} \hat{Q}_{i+1} A \hat{Q}_i$   
 $i \to \infty$   
 $A: = 0, A \hat{Q}_i$   
 $A \hat{Q}_i = \hat{Q}_i A \hat{Q}_i$   
 $A \hat{Q}_i = \hat{Q}_i A \hat{Q}_i = \hat{Q}_i \hat{Q} \hat{Q}_i$   
 $\hat{Q}_{i+1} \hat{R}_{i+1} = \hat{Q} \hat{Q}_i R_i$   
 $\hat{Q}_{i+1} \hat{R}_{i+1} = \hat{R}_i \hat{Q}_i R_i$   
 $\hat{Q}_{i+1} \hat{R}_{i+1} = \hat{R}_i \hat{Q}_i \hat{Q}_i$   
 $= R_i \hat{Q}_i$ 

## **QR** Iteration with Shift

Describe QR iteration with shifting

LO: R.J. OR(A: - J.I)

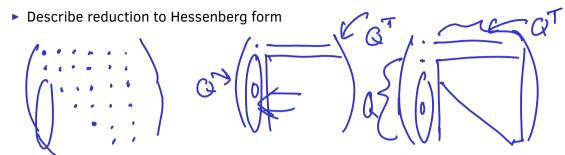
$$A_{i_{i_1}} = R_i Q_i + \sigma_i I$$

Discuss how shift can be selected

$$\sigma_i : (A_i)_{nn}$$

$$A_{i+1} = (A_i - \sigma_i) (A_i - \sigma_i)$$

## Hessenberg and Tridiagonal Form



Describe reduction to tridiagonal form in symmetric case

### **QR** Iteration Complexity

Compare complexity of QR iteration for various matrices

with Lidiagonal form BR. Lochen cuche O(n) per thether Keyescherg Form O(n2) per Arshon Cerved fimmer O(n3) per , heathor