

# CS 450: Numerical Analysis<sup>1</sup>

## Linear Systems

University of Illinois at Urbana-Champaign

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<sup>1</sup> *These slides have been drafted by Edgar Solomonik as lecture templates and supplementary material for the book “Scientific Computing: An Introductory Survey” by Michael T. Heath ([slides](#)).*

# Vector Norms

- ▶ **Properties of vector norms**
- ▶ **A norm is uniquely defined by its unit sphere:**
- ▶  **$p$ -norms**

# Inner-Product Spaces

- **Properties of inner-product spaces:** Inner products  $\langle x, y \rangle$  must satisfy

$$\langle x, x \rangle \geq 0$$

$$\langle x, x \rangle = 0 \quad \Leftrightarrow \quad x = \mathbf{0}$$

$$\langle x, y \rangle = \langle y, x \rangle$$

$$\langle x, y + z \rangle = \langle x, y \rangle + \langle x, z \rangle$$

$$\langle \alpha x, y \rangle = \alpha \langle x, y \rangle$$

- **Inner-product-based vector norms**

# Matrix Norms

► Properties of matrix norms:

$$||\mathbf{A}|| \geq 0$$

$$||\mathbf{A}|| = 0 \Leftrightarrow \mathbf{A} = \mathbf{0}$$

$$||\alpha\mathbf{A}|| = |\alpha| \cdot ||\mathbf{A}||$$

$$||\mathbf{A} + \mathbf{B}|| \leq ||\mathbf{A}|| + ||\mathbf{B}|| \quad (\textit{triangle inequality})$$

► Frobenius norm:

► Operator/induced/subordinate matrix norms:

# Induced Matrix Norms

- ▶ **Interpreting induced matrix norms:**

- ▶ **General induced matrix norms:**

# Matrix Condition Number

*Demo: Conditioning of 2x2 Matrices*

*Demo: Condition number visualized*

- ▶ **Definition:**  $\kappa(\mathbf{A}) = \|\mathbf{A}\| \cdot \|\mathbf{A}^{-1}\|$  is the ratio between the shortest/longest distances from the unit-ball center to any point on the surface.
- ▶ **Intuitive derivation:**

$$\kappa(\mathbf{A}) = \max_{\text{inputs}} \max_{\text{perturbations in input}} \left| \frac{\text{relative perturbation in output}}{\text{relative perturbation in input}} \right|$$

since a matrix is a linear operator, we can decouple its action on the input  $x$  and the perturbation  $\delta x$  since  $\mathbf{A}(x + \delta x) = \mathbf{A}x + \mathbf{A}\delta x$ , so

$$\kappa(\mathbf{A}) = \left| \frac{\overbrace{\max_{\text{perturbations in input}} \text{relative perturbation growth}}^{\|\mathbf{A}\|}}{\underbrace{\max_{\text{inputs}} \text{relative input reduction}}_{1/\|\mathbf{A}^{-1}\|}} \right|$$

## Matrix Conditioning

- ▶ The matrix condition number  $\kappa(\mathbf{A})$  is the ratio between the max and min distance from the surface to the center of the unit ball transformed by  $\kappa(\mathbf{A})$ :
- ▶ The matrix condition number bounds the worst-case amplification of error in a matrix-vector product:

# Norms and Conditioning of Orthogonal Matrices

- ▶ **Orthogonal matrices:**
- ▶ **Norm and condition number of orthogonal matrices:**



# Singular Value Decomposition

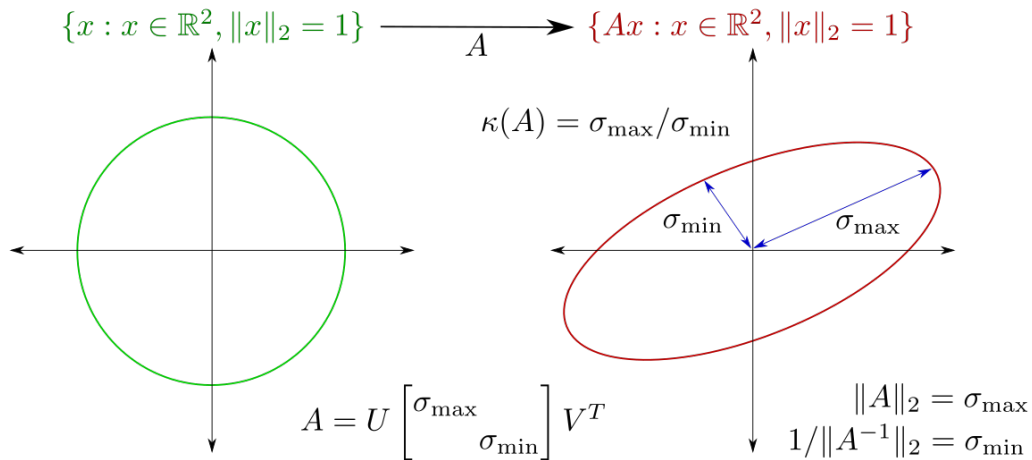
- ▶ **The singular value decomposition (SVD):**

# Norms and Conditioning via SVD

*Activity: Singular Value Decomposition and Norms*

- ▶ **Norm and condition number in terms of singular values:**

# Visualization of Matrix Conditioning



# Conditioning of Linear Systems

- ▶ Lets now return to formally deriving the conditioning of solving  $Ax = b$ :

## Conditioning of Linear Systems II

- ▶ Consider perturbations to the input coefficients  $\hat{A} = A + \delta A$ :

# Solving Basic Linear Systems

- ▶ Solve  $Dx = b$  if  $D$  is diagonal
- ▶ Solve  $Qx = b$  if  $Q$  is orthogonal
- ▶ Given SVD  $A = U\Sigma V^T$ , solve  $Ax = b$

# Solving Triangular Systems

- ▶  $Lx = b$  if  $L$  is lower-triangular is solved by forward substitution:

$$\begin{array}{rcl} l_{11}x_1 & = & b_1 \qquad x_1 = \\ l_{21}x_1 + l_{22}x_2 & = & b_2 \quad \Rightarrow \quad x_2 = \\ l_{31}x_1 + l_{32}x_2 + l_{33}x_3 & = & b_3 \qquad x_3 = \\ \vdots & & \vdots \end{array}$$

- ▶ Algorithm can also be formulated recursively by blocks:

# Solving Triangular Systems

- ▶ **Existence of solution to  $Lx = b$ :**
- ▶ **Uniqueness of solution:**
- ▶ **Computational complexity of forward/backward substitution:**



## Properties of Triangular Matrices

►  $Z = XY$  is lower triangular if  $X$  and  $Y$  are both lower triangular:

►  $L^{-1}$  is lower triangular if it exists:

## LU Factorization

- ▶ An **LU factorization** consists of a unit-diagonal lower-triangular **factor**  $L$  and upper-triangular factor  $U$  such that  $A = LU$ :
  
  
  
  
  
  
  
  
  
  
- ▶ Given an LU factorization of  $A$ , we can solve the linear system  $Ax = b$ :

# Gaussian Elimination Algorithm

- ▶ Algorithm for factorization is derived from equations given by  $A = LU$ :
- ▶ The computational complexity of LU is  $O(n^3)$ :

## Existence of LU Factorization

- ▶ **The LU factorization may not exist:** Consider matrix  $\begin{bmatrix} 3 & 2 \\ 6 & 4 \\ 0 & 3 \end{bmatrix}$ .

- ▶ **Permutation of rows enables us to transform the matrix so the LU factorization does exist:**

# Gaussian Elimination with Partial Pivoting

- ▶ **Partial pivoting** permutes rows to make divisor  $u_{ii}$  is maximal at each step:
- ▶ A row permutation corresponds to an application of a **row permutation matrix**  $P_{jk} = I - (e_j - e_k)(e_j - e_k)^T$ :

## Partial Pivoting Example

- ▶ Lets consider again the matrix  $A = \begin{bmatrix} 3 & 2 \\ 6 & 4 \\ 0 & 3 \end{bmatrix}$ .

## Complete Pivoting

- ▶ ***Complete pivoting*** permutes rows and columns to make divisor  $u_{ii}$  is maximal at each step:
- ▶ Complete pivoting is noticeably more expensive than partial pivoting:

## Round-off Error in LU

► Lets consider factorization of  $\begin{bmatrix} \epsilon & 1 \\ 1 & 1 \end{bmatrix}$  where  $\epsilon < \epsilon_{\text{mach}}$ :

► Permuting the rows of  $A$  in partial pivoting gives  $PA = \begin{bmatrix} 1 & 1 \\ \epsilon & 1 \end{bmatrix}$



## Error Analysis of LU

- ▶ The main source of round-off error in LU is in the computation of the Schur complement:
- ▶ When computed in floating point, absolute backward error  $\delta A$  in LU (so  $\hat{L}\hat{U} = A + \delta A$ ) is  $|\delta a_{ij}| \leq \epsilon_{\text{mach}}(|\hat{L}| \cdot |\hat{U}|)_{ij}$

## Helpful Matrix Properties

- ▶ Matrix is ***diagonally dominant***, so  $\sum_{i \neq j} |a_{ij}| \leq |a_{ii}|$ :
- ▶ Matrix is ***symmetric positive definite (SPD)***, so  $\forall_{x \neq 0}, x^T A x > 0$ :
- ▶ Matrix is symmetric but indefinite:
- ▶ Matrix is ***banded***,  $a_{ij} = 0$  if  $|i - j| > b$ :

# Solving Many Linear Systems

**Demo:** Sherman-Morrison

**Activity:** Sherman-Morrison-Woodbury Formula

- ▶ Suppose we have computed  $A = LU$  and want to solve  $AX = B$  where  $B$  is  $n \times k$  with  $k < n$ :
- ▶ Suppose we have computed  $A = LU$  and now want to solve a perturbed system  $(A - uv^T)x = b$ :  
Can use the *Sherman-Morrison-Woodbury* formula

$$(A - uv^T)^{-1} = A^{-1} + \frac{A^{-1}uv^T A^{-1}}{1 - v^T A^{-1}u}$$