

CS 450: Numerical Analysis¹

Linear Systems

$$\underline{A} \underline{x} = \underline{b}$$

?!

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¹ These slides have been drafted by Edgar Solomonik as lecture templates and supplementary material for the book “Scientific Computing: An Introductory Survey” by Michael T. Heath ([slides](#)).

Vector Norms

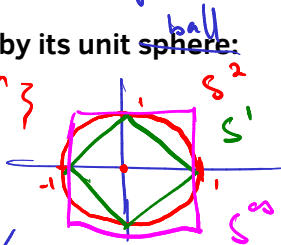
$$x \in \mathbb{R}^n$$

► Properties of vector norms

$$\begin{cases} \|x\| \geq 0, \quad \|x\| = 0 \text{ if } x = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix} \\ \|\alpha x\| = |\alpha| \|x\|, \quad \|x + y\| \leq \|x\| + \|y\| \end{cases}$$

► A norm is uniquely defined by its unit ~~sphere~~ ^{ball}:

$$S^p = \{x : \|x\|_p = 1, x \in \mathbb{R}^n\}$$

► p -norms

$$\|x\|_p = \left(\sum_{i=1}^n |x_i|^p \right)^{1/p}$$

$$\begin{aligned} \|x\|_2 &= \sqrt{\sum_{i=1}^n x_i^2} \\ \|x\|_1 &= \sum_{i=1}^n |x_i| \\ \|x\|_\infty &= \max_i |x_i| \end{aligned}$$

Inner-Product Spaces

- **Properties of inner-product spaces:** Inner products $\langle x, y \rangle$ must satisfy

$$\|x\|_2 = \sqrt{\langle x, x \rangle}$$

$$\langle x, x \rangle \geq 0$$

$$\langle x, x \rangle = 0 \Leftrightarrow x = 0$$

$$\langle x, y \rangle = \langle y, x \rangle \leftarrow \text{commutativity}$$

$$\langle x, y + z \rangle = \langle x, y \rangle + \langle x, z \rangle$$

$$\langle \alpha x, y \rangle = \alpha \langle x, y \rangle$$

$$= \sum_i x_i y_i$$

\leftarrow distributivity

- **Inner-product-based vector norms**

$$|\langle x, y \rangle| \leq \sqrt{\frac{\langle x, x \rangle \langle y, y \rangle}{\|x\|_2^2}}$$

$$|\langle x, y \rangle| \leq \|x\|_2 \|y\|_2$$

Matrix Norms

► Properties of matrix norms:

$$\|A\| \geq 0$$

$$\|A\| = 0 \Leftrightarrow A = 0$$

$$\|\alpha A\| = |\alpha| \cdot \|A\|$$

$$\|A + B\| \leq \|A\| + \|B\| \quad (\text{triangle inequality})$$

$$\text{vec}([a_1 \dots a_n]) \rightarrow$$

$$\begin{bmatrix} a_1 \\ a_2 \\ \vdots \\ a_n \end{bmatrix}$$

► Frobenius norm:

$$\|A\|_F = \sqrt{\sum_{i=1}^n \sum_{j=1}^n a_{ij}^2} = \|\text{vec}(A)\|_2$$

► Operator/induced/subordinate matrix norms:

$$A \rightarrow f(x) = Ax$$

$$\|A\|_q = \max_{\substack{x \in \mathbb{R}^n \\ x \neq 0}} \|Ax\|_q$$

$$\frac{\|Ax\|_q}{\|x\|_q} = \max_{x \in \mathbb{R}^n} \|Ax\|_q$$

$$\|x\|_q$$

Induced Matrix Norms

► Interpreting induced matrix norms:

$\|A\|_p \rightarrow$ max. magn. ampl. factor for the norm
of Ax relative to x over all x

$$\min_{\substack{x \in \mathbb{R}^n \\ \|x\|_p = 1}} \|Ax\|_p = \frac{1}{\|A^{-1}\|_p} \quad \left| \quad y = Ax, \quad A^{-1}y = x \right.$$

► General induced matrix norms:

Matrix Condition Number

Demo: Conditioning of 2x2 Matrices

Demo: Condition number visualized

- **Definition:** $\kappa(\mathbf{A}) = \|\mathbf{A}\| \cdot \|\mathbf{A}^{-1}\|$ is the ratio between the shortest/longest distances from the unit-ball center to any point on the surface.

- **Intuitive derivation:**

$$\kappa(\mathbf{A}) = \max_{\text{inputs}} \max_{\text{perturbations in input}} \left| \frac{\text{relative perturbation in output}}{\text{relative perturbation in input}} \right|$$

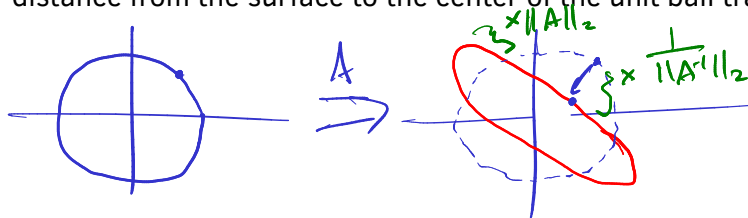
since a matrix is a linear operator, we can decouple its action on the input x and the perturbation δx since $\mathbf{A}(x + \delta x) = \mathbf{A}x + \mathbf{A}\delta x$, so

$$\begin{aligned} \|\mathbf{A}\delta x\| &\leq \|\mathbf{A}\| \|\delta x\| \\ \|\mathbf{A}x\| &\geq \frac{1}{\|\mathbf{A}^{-1}\|} \|x\| \end{aligned}$$

$$\kappa(\mathbf{A}) = \frac{\overbrace{\max_{\text{perturbations in input}} \text{relative perturbation growth}}^{\|\mathbf{A}\|}}{\underbrace{\max_{\text{inputs}} \text{relative input reduction}}_{1/\|\mathbf{A}^{-1}\|}}$$

Matrix Conditioning

- ▶ The matrix condition number $\kappa(\mathbf{A})$ is the ratio between the max and min distance from the surface to the center of the unit ball transformed by $\kappa(\mathbf{A})$:



- ▶ The matrix condition number bounds the worst-case amplification of error in a matrix-vector product:

$$\frac{\|\delta y\|}{\|y\|} \leq \frac{\|\delta x\|}{\|x\|}$$

$\|A\| \|A^{-1}\|_2$

$$y + \delta y = A(x + \delta x) = Ax + A\delta x$$

$$\delta y = A\delta x$$

$$\|\delta y\|_2 \leq \|A\|_2 \|\delta x\|_2$$

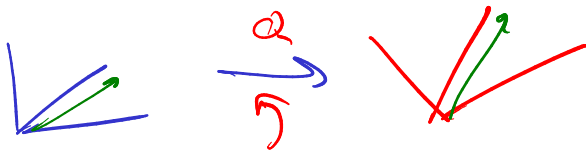
$$\|y\|_2 \geq \frac{1}{\|A^{-1}\|_2} \|x\|_2$$

Norms and Conditioning of Orthogonal Matrices

- **Orthogonal matrices:** $Q^T = Q^{-1}$ $Q^T Q = I$ $Q Q^T = I$
 $Q \in \mathbb{R}^{n \times n}$

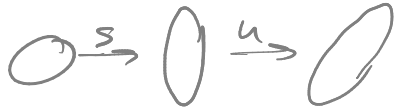
- **Norm and condition number of orthogonal matrices:**

$$\|Q\|_2 = 1 \quad \|Qx\|_2 = \|x\|_2$$



$$\kappa(Q) = \|Q\|_2 \|Q^{-1}\|_2 \stackrel{Q^T}{=} 1$$

Singular Value Decomposition

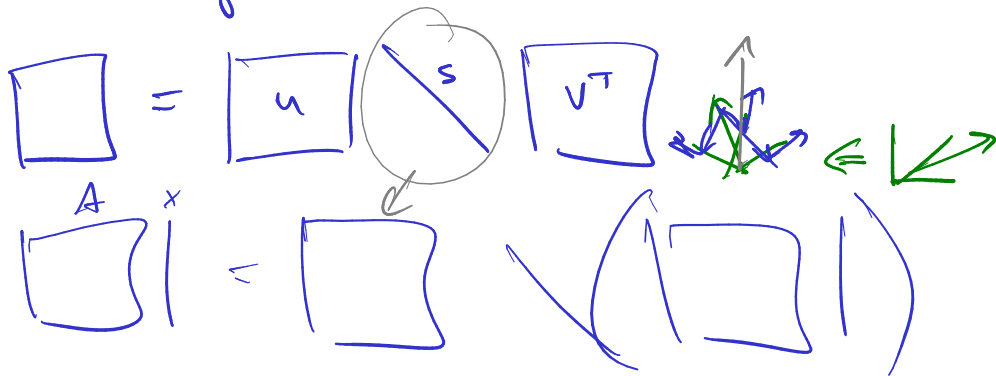


- The singular value decomposition (SVD):

$$A = U S V^T$$

↑ ↑
orthogonal

diagonal and nonnegative
(positive if A is full rank)



Norms and Conditioning via SVD

Activity: Singular Value Decomposition and Norms

- Norm and condition number in terms of singular values:

$$S = \begin{bmatrix} \sigma_{\max} & & \\ & \ddots & \\ & & \sigma_{\min} \end{bmatrix},$$

$$\|S\|_2 =$$

$$\Rightarrow \|A\| = \cancel{\|U\|} \cancel{\|S\|} \cancel{\|V\|}$$

$$A = U S V^T$$

$$S = \begin{bmatrix} \sigma_1 & & \\ & \ddots & \\ & & \sigma_n \end{bmatrix} \quad \text{singular values}$$

$$U^T A V = S$$

$$\max_{\|x\|_2=1} \|Sx\|_2 = \max_{\|x\|_2=1}$$

$$\sqrt{\sum_i \sigma_i^2 x_i^2}$$

$$\|S\| \geq \cancel{\|U\|} \cdot \cancel{\|A\|} \cdot \cancel{\|V\|}$$

$$\|S\|_2 = \sigma_{\max}$$

$$\|S^{-1}\|_2 = \frac{1}{\sigma_{\min}}$$

$$\|A\| \|A^{-1}\| = \|S\| \|S^{-1}\| = \sigma_{\max} / \sigma_{\min}$$

Visualization of Matrix Conditioning

$$(A^T A)^T = A^T A$$

$$\frac{x^T A^T (A x)}{x^T x} \geq 0$$

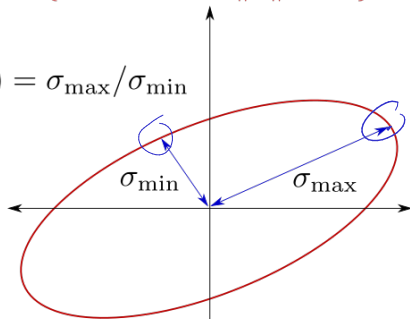
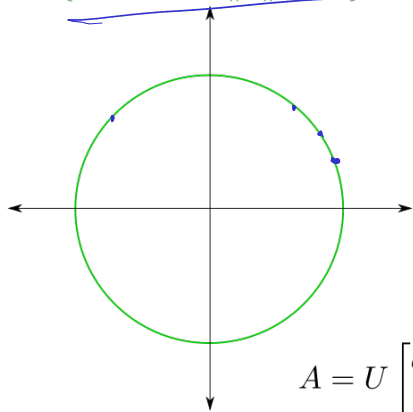
$$\frac{y^T y}{x^T x} \geq 0$$

$$\{x : x \in \mathbb{R}^2, \|x\|_2 = 1\}$$

A

$$\{Ax : x \in \mathbb{R}^2, \|x\|_2 = 1\}$$

$$\kappa(A) = \sigma_{\max} / \sigma_{\min}$$



$$A = U \begin{bmatrix} \sigma_{\max} & \\ & \sigma_{\min} \end{bmatrix} V^T$$

$$\|A\|_2 = \sigma_{\max}$$

$$1/\|A^{-1}\|_2 = \sigma_{\min}$$

$$A^{-1} = V \begin{bmatrix} \frac{1}{\sigma_{\max}} & \\ & \frac{1}{\sigma_{\min}} \end{bmatrix} U^T$$

Conditioning of Linear Systems

- ▶ Lets now return to formally deriving the conditioning of solving $Ax = b$:

Conditioning of Linear Systems II

- ▶ Consider perturbations to the input coefficients $\hat{A} = A + \delta A$:

Solving Basic Linear Systems

- ▶ Solve $Dx = b$ if D is diagonal
- ▶ Solve $Qx = b$ if Q is orthogonal
- ▶ Given SVD $A = U\Sigma V^T$, solve $Ax = b$

Solving Triangular Systems

- ▶ $Lx = b$ if L is lower-triangular is solved by forward substitution:

$$\begin{array}{rcl} l_{11}x_1 & = & b_1 \\ l_{21}x_1 + l_{22}x_2 & = & b_2 \\ l_{31}x_1 + l_{32}x_2 + l_{33}x_3 & = & b_3 \\ \vdots & & \vdots \end{array} \quad \Rightarrow \quad \begin{array}{l} x_1 = \\ x_2 = \\ x_3 = \\ \vdots \end{array}$$

- ▶ Algorithm can also be formulated recursively by blocks:

Solving Triangular Systems

- ▶ **Existence of solution to $Lx = b$:**
- ▶ **Uniqueness of solution:**
- ▶ **Computational complexity of forward/backward substitution:**

Properties of Triangular Matrices

► $Z = XY$ is lower triangular if X and Y are both lower triangular:

► L^{-1} is lower triangular if it exists:

LU Factorization

- ▶ An **LU factorization** consists of a unit-diagonal lower-triangular **factor** L and upper-triangular factor U such that $A = LU$:

- ▶ Given an LU factorization of A , we can solve the linear system $Ax = b$:

Gaussian Elimination Algorithm

- ▶ Algorithm for factorization is derived from equations given by $A = LU$:
- ▶ The computational complexity of LU is $O(n^3)$:

Existence of LU Factorization

- ▶ **The LU factorization may not exist:** Consider matrix $\begin{bmatrix} 3 & 2 \\ 6 & 4 \\ 0 & 3 \end{bmatrix}$.

- ▶ **Permutation of rows enables us to transform the matrix so the LU factorization does exist:**

Gaussian Elimination with Partial Pivoting

► **Partial pivoting** permutes rows to make divisor u_{ii} is maximal at each step:

► A row permutation corresponds to an application of a **row permutation matrix** $P_{jk} = I - (e_j - e_k)(e_j - e_k)^T$:

Partial Pivoting Example

- ▶ Lets consider again the matrix $A = \begin{bmatrix} 3 & 2 \\ 6 & 4 \\ 0 & 3 \end{bmatrix}$.

Complete Pivoting

- ▶ ***Complete pivoting*** permutes rows and columns to make divisor u_{ii} is maximal at each step:

- ▶ Complete pivoting is noticeably more expensive than partial pivoting:

Round-off Error in LU

► Lets consider factorization of $\begin{bmatrix} \epsilon & 1 \\ 1 & 1 \end{bmatrix}$ where $\epsilon < \epsilon_{\text{mach}}$:

► Permuting the rows of A in partial pivoting gives $PA = \begin{bmatrix} 1 & 1 \\ \epsilon & 1 \end{bmatrix}$

Error Analysis of LU

- ▶ The main source of round-off error in LU is in the computation of the Schur complement:
- ▶ When computed in floating point, absolute backward error δA in LU (so $\hat{L}\hat{U} = A + \delta A$) is $|\delta a_{ij}| \leq \epsilon_{\text{mach}}(|\hat{L}| \cdot |\hat{U}|)_{ij}$

Helpful Matrix Properties

- ▶ Matrix is ***diagonally dominant***, so $\sum_{i \neq j} |a_{ij}| \leq |a_{ii}|$:
- ▶ Matrix is ***symmetric positive definite (SPD)***, so $\forall_{x \neq 0}, x^T A x > 0$:
- ▶ Matrix is symmetric but indefinite:
- ▶ Matrix is ***banded***, $a_{ij} = 0$ if $|i - j| > b$:

Solving Many Linear Systems

Demo: Sherman-Morrison

Activity: Sherman-Morrison-Woodbury Formula

- ▶ Suppose we have computed $A = LU$ and want to solve $AX = B$ where B is $n \times k$ with $k < n$:
- ▶ Suppose we have computed $A = LU$ and now want to solve a perturbed system $(A - uv^T)x = b$:
Can use the *Sherman-Morrison-Woodbury* formula

$$(A - uv^T)^{-1} = A^{-1} + \frac{A^{-1}uv^T A^{-1}}{1 - v^T A^{-1}u}$$