

# CS 450: Numerical Analysis<sup>1</sup>

## Linear Least Squares

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<sup>1</sup>*These slides have been drafted by Edgar Solomonik as lecture templates and supplementary material for the book “Scientific Computing: An Introductory Survey” by Michael T. Heath ([slides](#)).*

## Linear Least Squares

- ▶ Find  $\mathbf{x}^* = \operatorname{argmin}_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2$  where  $\mathbf{A} \in \mathbb{R}^{m \times n}$ .
  
- ▶ Given the SVD  $\mathbf{A} = \mathbf{U}\mathbf{\Sigma}\mathbf{V}^T$  we have  $\mathbf{x}^* = \mathbf{V}\mathbf{\Sigma}^\dagger\mathbf{U}^T\mathbf{b}$ , where  $\mathbf{\Sigma}^\dagger$  contains the reciprocal of all nonzeros in  $\mathbf{\Sigma}$ :

## Conditioning of Linear Least Squares

- ▶ Consider fitting a line to a collection of points, then perturbing the points:
  
  
  
  
  
  
  
  
  
  
- ▶ LLS is ill-posed for any  $A$ , unless we consider solving for a particular  $b$

# Normal Equations

*Demo: Normal equations vs Pseudoinverse*

*Demo: Issues with the normal equations*

- ▶ *Normal equations* are given by solving  $A^T A x = A^T b$ :
  
  
  
  
  
  
  
  
  
  
- ▶ However, solving the normal equations is a more ill-conditioned problem than the original least squares algorithm

## Solving the Normal Equations

- ▶ If  $A$  is full-rank, then  $A^T A$  is symmetric positive definite (SPD):
  
  
  
  
  
  
  
  
  
  
  
- ▶ Since  $A^T A$  is SPD we can use Cholesky factorization, to factorize it and solve linear systems:

LLS

$$\begin{array}{c}
 A \\
 \boxed{\phantom{A}}
 \end{array}
 \begin{array}{c}
 ? \\
 x \\
 |
 \end{array}
 \approx
 \begin{array}{c}
 b \\
 |
 \end{array}$$

$$\underbrace{\|Ax - b\|_2}_{\text{residual}}$$

$$\begin{array}{c}
 \downarrow \\
 \text{span}(A) = \text{span}(Q)
 \end{array}
 \quad
 \begin{array}{c}
 Q^T Q = I \\
 \leftarrow (I - QQ^T)b = Ax - b
 \end{array}$$

QR factorization:  $A = QR$

$$\boxed{\phantom{A}} = \boxed{\phantom{Q}} \boxed{\phantom{R}} = \boxed{\phantom{Q}} \begin{array}{c} \triangle \\ \phantom{R} \end{array}$$

## QR Factorization

- ▶ If  $A$  is full-rank there exists an orthogonal matrix  $Q$  and a unique upper-triangular matrix  $R$  with a positive diagonal such that  $A = QR$
  
- ▶ A reduced QR factorization (unique part of general QR) is defined so that  $Q \in \mathbb{R}^{m \times n}$  has orthonormal columns and  $R$  is square and upper-triangular

- ▶ We can solve the normal equations (and consequently the linear least squares problem) via reduced QR as follows

$$A = QR \quad \left| \quad \begin{array}{l} Q^T A x = Q^T b \\ R x = Q^T b \end{array} \quad \begin{array}{l} A \in \mathbb{R}^{m \times n} \\ \mathcal{O}(mn) \end{array}$$

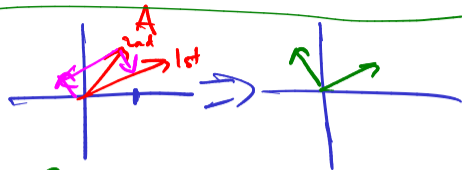
# Gram-Schmidt Orthogonalization

Demo: Gram-Schmidt-The Movie

Demo: Gram-Schmidt and Modified Gram-Schmidt

► Classical Gram-Schmidt process for QR:

$$P = u u^T \quad \left| \begin{array}{l} P v = u \cdot \langle u, v \rangle \\ (I - u u^T) v = w \\ w \perp u = 0 \end{array} \right.$$



► Modified Gram-Schmidt process for QR:

$$a_j^{(i)} = a_j \quad a_j^{(i)} \rightarrow a_j^{(i-1)} - \sum_{k=1}^{i-1} q_k \langle q_k, a_j^{(i-1)} \rangle$$

$$\left[ \begin{array}{c} | \\ | \\ | \\ | \end{array} \right] \xrightarrow{A} \left[ \begin{array}{c} | \\ | \\ | \\ | \end{array} \right] \xrightarrow{Q} \left[ \begin{array}{c} | \\ | \\ | \\ | \end{array} \right] \left[ \begin{array}{c} \cdot \\ \vdots \\ \cdot \end{array} \right]$$



$$u_i = a_i - \sum_{j < i} q_j \langle q_j, a_i \rangle$$

$$u_i = \begin{bmatrix} I & P_{j < i} \\ P_1 & P_2 & P_3 \end{bmatrix} a_i$$

$$P_j = I - q_j q_j^T$$



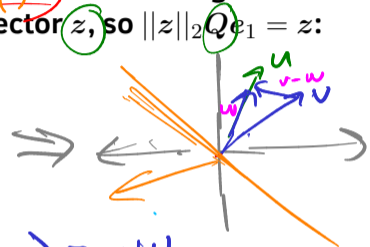
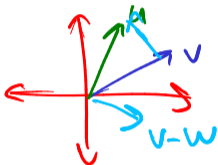
# Householder QR Factorization

- ▶ A Householder transformation  $Q = I - 2uu^T$  is an orthogonal matrix defined to annihilate entries of a given vector  $z$ , so  $\|z\|_2 Qe_1 = z$ :

$I - 2uu^T$  ← MGS/CGS

$Qz = \|z\|_2 e_1$

$uu^T v = w$



GS  $\|v-w\| \leq \|v\|$

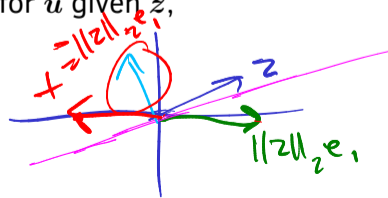
Householder  $\|v-2w\| = \|v\|$

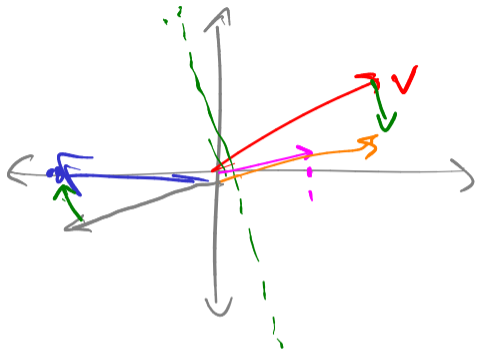
$uu^T (v-2w) = -w$

- ▶ Imposing this form on  $Q$  leaves exactly two choices for  $u$  given  $z$ ,

$Q = I - 2uu^T$   
 $(I - 2uu^T)z = \|z\|_2 e_1$

$u = \frac{z \pm \|z\|_2 e_1}{\|z \pm \|z\|_2 e_1\|_2}$

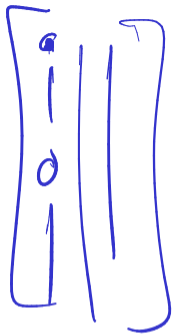




Q



=



## Applying Householder Transformations

- ▶ The product  $x = Qw$  can be computed using  $O(n)$  operations if  $Q$  is a Householder transformation
  
- ▶ Householder transformations are also called *reflectors* because their application reflects a vector along a hyperplane (changes sign of component of  $w$  that is parallel to  $u$ )

## Givens Rotations

- ▶ Householder reflectors reflect vectors, Givens rotations rotate them

- ▶ Givens rotations are defined by orthogonal matrices of the form  $\begin{bmatrix} c & s \\ -s & c \end{bmatrix}$



## Rank-Deficient Least Squares

- ▶ Suppose we want to solve a linear system or least squares problem with a (nearly) rank deficient matrix  $A$
  
  
  
  
  
  
  
  
  
  
- ▶ Rank-deficient least squares problems seek a minimizer  $x$  of  $\|Ax - b\|_2$  of minimal norm  $\|x\|_2$

## Truncated SVD

- ▶ After floating-point rounding, rank-deficient matrices typically regain full-rank but have nonzero singular values on the order of  $\epsilon_{\text{mach}}\sigma_{\text{max}}$
  
  
  
  
  
  
  
  
  
  
- ▶ By the *Eckart-Young-Mirsky theorem*, truncated SVD also provides the best low-rank approximation of a matrix (in 2-norm and Frobenius norm)

