

CS 450: Numerical Analysis¹

Fast Fourier Transform

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¹These slides have been drafted by Edgar Solomonik as lecture templates and supplementary material for the book “Scientific Computing: An Introductory Survey” by Michael T. Heath ([slides](#)).

Multigrid

- ▶ Consider the Galerkin approximation with linear finite elements to the Poisson equation $u'' = f(t)$ with boundary conditions $u(a) = u(b) = 0$:

$$\phi_i^{(h)}(t) = \begin{cases} (t - t_{i-1})/h & : t \in [t_{i-1}, t_i] \\ (t_{i+1} - t)/h & : t \in [t_i, t_{i+1}] \\ 0 & : \text{otherwise} \end{cases}$$

where $t_0 = t_1 = a$ and $t_{n+1} = t_n = b$.

Coarse Grid Matrix

- ▶ Multigrid restricts the residual equation on the fine grid $\mathbf{A}^{(h)}\mathbf{x} = \mathbf{r}^{(h)}$ to the coarse grid:

Restricting the Residual Equation

- ▶ Given the fine-grid residual $\mathbf{r}^{(h)}$, we seek to use the coarse grid to approximate $\mathbf{x}^{(h)}$ so that $\mathbf{A}\mathbf{x}^{(h)} \approx \mathbf{r}^{(h)}$

Discrete Fourier Transform

- ▶ The solutions to hyperbolic PDEs like Poisson are wave-like and take on simple representations in the frequency basis, both for continuous and discretized equations. We define the *discrete Fourier transform* using

$$\omega_{(n)} = \cos(2\pi/n) - i \sin(2\pi/n) = e^{-2\pi i/n}.$$

Fast Fourier Transform (FFT)

- ▶ Consider $\mathbf{b} = \mathbf{F}\mathbf{a}$, we have

$$\forall j \in [0, n-1] \quad b_j = \sum_{k=0}^{n-1} \omega_{(n)}^{jk} a_k,$$

the FFT computes this recursively via 2 FFTs of dimension $n/2$, using

$$\omega_{(n/2)} = \omega_{(n)}^2,$$

Fast Fourier Transform Derivation

- ▶ The FFT leverages similarity between the first and second half of the output,

$$b_j = \underbrace{\sum_{k=0}^{n/2-1} \omega_{(n/2)}^{jk} a_{2k}}_{u_j} + \omega_{(n)}^j \underbrace{\sum_{k=0}^{n/2-1} \omega_{(n/2)}^{jk} a_{2k+1}}_{v_j}$$

corresponds closely to the entry shifted by $n/2$,

$$b_{j+n/2} = \sum_{k=0}^{n/2-1} \omega_{(n/2)}^{(j+n/2)k} a_{2k} + \omega_{(n)}^{j+n/2} \sum_{k=0}^{n/2-1} \omega_{(n/2)}^{(j+n/2)k} a_{2k+1}$$

FFT Algorithm Summary

- ▶ Let vectors u and v be two recursive FFTs, $\forall j \in [0, n/2 - 1]$

$$u_j = \sum_{k=0}^{n/2-1} \omega_{(n/2)}^{jk} a_{2k}, \quad v_j = \sum_{k=0}^{n/2-1} \omega_{(n/2)}^{jk} a_{2k+1}$$

- ▶ The FFT has $O(n \log n)$ cost complexity:

Convolution via DFT

- ▶ The Fourier transform method for computing a convolution is given by

$$c_k = \frac{1}{n} \sum_s \omega_{(n)}^{-ks} \left(\sum_j \omega_{(n)}^{sj} a_j \right) \left(\sum_t \omega_{(n)}^{st} b_t \right)$$

Solving Numerical PDEs with the FFT

- ▶ 1D finite-difference schemes on a regular grid correspond to convolutions:
- ▶ For the 1D Poisson model problem, the eigenvectors of T corresponds to the imaginary part of a minor of a $2(n + 1)$ -dimensional DFT matrix:
- ▶ Multidimensional Poisson can be handled with multidimensional FFT: