

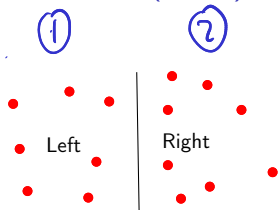
Ann:

Review:

Goals:

- Direct solve BSS mat.
→ HBSS mat.
- Butterfly
- PDE solvers / compact op.
theory

(Recursive) Coordinate Bisection (RCB)



	①	②
①	full rank	low rank
②	low rank	full rank

Block-separable matrices

Particles:



$$A = \begin{bmatrix} D_1 & A_{12} & A_{13} & A_{14} \\ A_{21} & D_2 & A_{23} & A_{24} \\ A_{31} & A_{32} & D_3 & A_{34} \\ A_{41} & A_{42} & A_{43} & D_4 \end{bmatrix}$$

D_i : Full rank
 if sym: $A_{ij} = A_{ji}^T$

where A_{ij} has low rank: How to capture rank structure?

$$A_{ij} \approx P_i \underbrace{(A_{ij})_{\mathbb{I}, \mathbb{J}}}_{\tilde{A}_{ij}} \Pi_j$$

\uparrow \uparrow \uparrow
 $O(n)$ \tilde{A}_{ij} $O(n)$

Optimistically:

$K \times K$

K : off-diag
 interaction rank

Proxy Recap

Saw: If A comes from a kernel for which Green's formula holds, then the same skeleton will work for all of space, for a given set of sources/targets. What would the resulting matrix look like?

Rank and Proxies



Unlike FMMs, partitions here do not include “buffer” zones of near elements. What are the consequences?

Idea:

- Add extra proxies from nearby particle groups
- Tolerate the increased rank,

Block-Separable Matrices

A *block-separable matrix* looks like this:

$$A_{i,j} = \begin{bmatrix} \square & \square \\ \square & \square \end{bmatrix} = \begin{bmatrix} \square \\ \square \end{bmatrix} \begin{bmatrix} \square & \square \end{bmatrix}$$

$$A \approx \begin{bmatrix} D_1 & P_1 \tilde{A}_{12} \Pi_2 & P_1 \tilde{A}_{13} \Pi_3 & P_1 \tilde{A}_{14} \Pi_4 \\ P_2 \tilde{A}_{21} \Pi_1 & D_2 & P_2 \tilde{A}_{23} \Pi_3 & P_2 \tilde{A}_{24} \Pi_4 \\ P_3 \tilde{A}_{31} \Pi_1 & P_3 \tilde{A}_{32} \Pi_2 & D_3 & P_3 \tilde{A}_{34} \Pi_4 \\ P_4 \tilde{A}_{41} \Pi_1 & P_4 \tilde{A}_{42} \Pi_2 & P_4 \tilde{A}_{43} \Pi_3 & D_4 \end{bmatrix}$$

$$C = \begin{pmatrix} 0 & \tilde{A}_{12} & \dots \\ \tilde{A}_{21} & & \\ & & \\ & & & 0 \end{pmatrix}$$

Here:

- ▶ \tilde{A}_{ij} smaller than A_{ij}
- ▶ D_i has full rank (not necessarily diagonal)
- ▶ P_i shared for entire row
- ▶ Π_i shared for entire column

For matrices: $D \overset{N \times N}{\times} + P \overset{N \times K}{\times} C \overset{K \times N}{\times} \Pi \overset{N \times N}{\times}$

Shape of Π_i :

$$\Pi_i = \begin{bmatrix} \square & \square & \square & \square \end{bmatrix}$$

$$\Pi = \begin{pmatrix} \Pi_1 \\ \vdots \\ \Pi_4 \end{pmatrix} \quad \tilde{X} = \Pi \times$$

Block-Separable Matrix: Questions

Q: Why is it called that?

$A = \begin{bmatrix} | & & | \\ \hline & & \\ \hline | & & | \end{bmatrix}$ b_i

$u(x,y) = \sum_i a_i(x) b_i(y)$
"separation of variables"

Q: How expensive is a matvec?

$O(N^{3/2})$ by analogy to single-level Barnes-Hut.

Q: How about a solve?

?

assmily: \sqrt{N} particles in a box/groups.

BSS Solve (I)

Separate out 'coarse' unknowns. Use the following notation:

$$B = \begin{bmatrix} 0 & P_1 \tilde{A}_{12} & P_1 \tilde{A}_{13} & P_1 \tilde{A}_{14} \\ P_2 \tilde{A}_{21} & 0 & P_2 \tilde{A}_{23} & P_2 \tilde{A}_{24} \\ P_3 \tilde{A}_{31} & P_3 \tilde{A}_{32} & 0 & P_3 \tilde{A}_{34} \\ P_4 \tilde{A}_{41} & P_4 \tilde{A}_{42} & P_4 \tilde{A}_{43} & 0 \end{bmatrix} \quad A_{\tilde{x}-b}$$

and

$$D = \begin{bmatrix} D_1 & & & \\ & D_2 & & \\ & & D_3 & \\ & & & D_4 \end{bmatrix}, \quad \Pi = \begin{bmatrix} \Pi_1 & & & \\ & \Pi_2 & & \\ & & \Pi_3 & \\ & & & \Pi_4 \end{bmatrix}.$$

$$A = D + B\Pi$$

$$\begin{pmatrix} D & B \\ -\Pi & Id \end{pmatrix} \begin{pmatrix} \tilde{x} \\ x_2 \end{pmatrix} = \begin{pmatrix} b \\ 0 \end{pmatrix} \Leftrightarrow \begin{matrix} D\tilde{x} + Bx_2 = b \\ -\Pi\tilde{x} + x_2 = 0 \end{matrix}$$

$\Leftrightarrow \Pi\tilde{x} = x_2$

BSS Solve (II)

Q: What are the matrix sizes? The vector lengths of \mathbf{x} and $\tilde{\mathbf{x}}$?

$$\Pi : \boxed{\quad}$$

Now work towards doing *just* a 'coarse' solve on $\tilde{\mathbf{x}}$, using the Schur complement:

$$\begin{pmatrix} D & B \\ -\Pi & Id \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} = \begin{pmatrix} b \\ 0 \end{pmatrix} \quad | \quad (\Pi D^{-1}) \cdot \rightsquigarrow \left(\Pi D^{-1} D \quad \Pi D^{-1} B \mid \Pi D^{-1} b \right)$$

$$\begin{pmatrix} 0 & Id + \Pi D^{-1} B \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} = \begin{pmatrix} \Pi D^{-1} b \end{pmatrix}$$

$$(Id + \Pi D^{-1} B) \tilde{\mathbf{x}} = \Pi D^{-1} b$$

BSS Solve (III)

Focus in on the second row:

$$* \quad \underbrace{(\text{Id} + \Pi D^{-1} B)}_{\substack{\text{on-dig} \\ \text{off-dig}}} \tilde{\mathbf{x}} = \Pi D^{-1} \mathbf{b} \quad \leftarrow$$

Every non-zero (i.e. off-diagonal) entry in $\Pi D^{-1} B$ looks like

$$\underbrace{\Pi_i D_i^{-1} P_i}_{\substack{\text{on-dig} \\ \text{off-dig}}} \tilde{A}_{ij}$$

Define a diagonal entry:

$$\tilde{A}_{ii} = (\Pi_i D_i^{-1} P_i)^{-1}$$

Idea: Hit \otimes from left with $\text{block-diag}_i(\tilde{A}_{ii})$

BSS Solve (IV)

Next, left-multiply $(\text{Id} + \Pi D^{-1} B)$ by block-diag $_i(\tilde{A}_{ii})$:

$$\begin{pmatrix} \hat{A}_{11} & & & \\ & \hat{A}_{22} & & \\ & & \hat{A}_{33} & \\ & & & \hat{A}_{44} \end{pmatrix} (\text{Id} + \Pi D^{-1} B) = \underbrace{\begin{pmatrix} \tilde{A}_{11} & \tilde{A}_{12} & & \\ \tilde{A}_{21} & & \ddots & \\ & & & \\ & & & \tilde{A}_{44} \end{pmatrix}}_{\tilde{A}}$$

$$(\text{Id} + \Pi D^{-1} B) \tilde{x} = \Pi D^{-1} b$$

$$\tilde{A} \tilde{x} = \text{block-diag}_i(\tilde{A}_{ii}) \Pi D^{-1} b$$

$$\Pi \tilde{x} = x$$

~~$$\tilde{x} = \Pi^{-1} x$$~~

$$\rightarrow D x + B x = b$$

BSS Solve: Summary



What have we achieved?

- ▶ Instead of solving a linear system of size

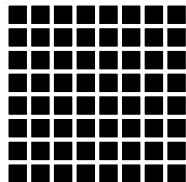
$$(N_{L0 \text{ boxes}} \cdot m) \times (N_{L0 \text{ boxes}} \cdot m)$$

we solve a linear system of size

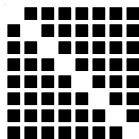
$$(N_{L0 \text{ boxes}} \cdot K) \times (N_{L0 \text{ boxes}} \cdot K),$$

which is cheaper by a factor of $(K/m)^3$.

- ▶ We are now only solving on the skeletons.



↓ compress



(Figure following G. Martinsson, drawn by A. Fikl)

Hierarchically Block-Separable

To get to $O(N)$, realize we can *recursively*

- ▶ group skeletons
- ▶ eliminate more variables.

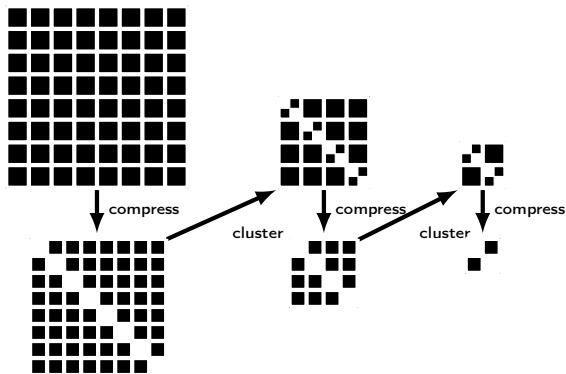
Where does this process start?



Demo: Skeletonization using Proxies (Hierarchical)

Hierarchically Block-Separable

In order to get $O(N)$ complexity, could we apply this procedure recursively?



(Figure following G. Martinsson, drawn by A. Fikl)

Hierarchically Block-Separable

- ▶ Using this hierarchical grouping gives us *Hierarchically Block-Separable (HBS)* matrices.
- ▶ If you have heard the word \mathcal{H} -matrix and \mathcal{H}^2 -matrix, the ideas are very similar. Differences:
 - ▶ \mathcal{H} -family matrices don't typically use the ID (instead often use *Adaptive Cross Approximation* or *ACA*)
 - ▶ \mathcal{H}^2 does target clustering (like FMM), \mathcal{H} does not (like Barnes-Hut)

Outline

Introduction

Dense Matrices and Computation

Tools for Low-Rank Linear Algebra

Rank and Smoothness

Near and Far: Separating out High-Rank Interactions

Ewald Summation

Barnes-Hut

Fast Multipole

Direct Solvers

The Butterfly Factorization

Outlook: Building a Fast PDE Solver

Going Infinite: Integral Operators and Functional Analysis

Singular Integrals and Potential Theory

Boundary Value Problems

Back from Infinity: Discretization

Computing Integrals: Approaches to Quadrature

Going General: More PDEs

Recap: Fast Fourier Transform

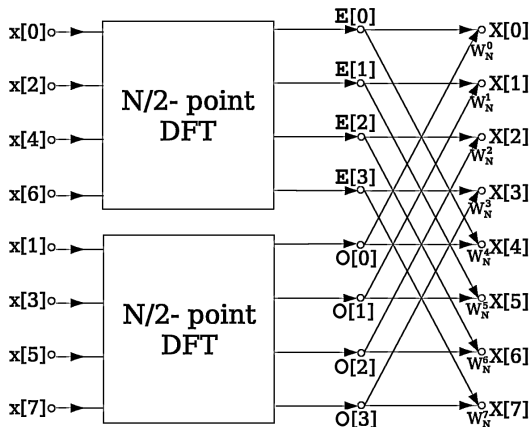
The *Discrete Fourier Transform (DFT)* is given by:

$$X_k = \sum_{n=0}^{N-1} x_n e^{-\frac{2\pi i}{N}nk} \quad (k = 0, \dots, N-1)$$

The foundation of the *Fast Fourier Transform (FFT)* is the factorization:

$$X_k = \underbrace{\sum_{m=0}^{N/2-1} x_{2m} e^{-\frac{2\pi i}{N/2}mk}}_{\text{DFT of even-indexed part of } x_n} + e^{-\frac{2\pi i}{N}k} \underbrace{\sum_{m=0}^{N/2-1} x_{2m+1} e^{-\frac{2\pi i}{N/2}mk}}_{\text{DFT of odd-indexed part of } x_n} .$$

FFT: Data Flow



(Figure credit: [Wikipedia](#))

Perhaps a little bit like a butterfly?

Fourier Transforms: A Different View

$$\mathcal{F}f(x) = \int e^{ixt} f(t) dt$$

Claim:

The [numerical] rank of the normalized Fourier transform with kernel e^{ixt} is bounded by a constant times γ , at any fixed precision ϵ .

(i.e. rank is proportional to the area of the rectangle swept out by x and t)

[[O'Neil et al. '10](#)]

Demo: Butterfly Factorization (Part I)

Recompression: Making use of Area-Bounded Rank

How do rectangular submatrices get expressed so as to reveal their constant rank?

