

Numerical Python

optimization

CS101 Lecture #20

Administrivia

- ❖ Homework #9 is due Friday, Dec. 9.
- ❖ Homework #10 is due Tuesday, Dec. 20.
- ❖ Midterm #2 is Monday, Dec. 19 from 7–10 p.m.

Warmup Question

Question #1

```
x = 'ABCD'  
z = 'XYZ'
```

```
for a in itertools.product( x,y ):  
    print( ' '.join( a ) )
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Which of the following is *not* printed?

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Brute-Force Search

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10	$86^{10} = 2.2 \times 10^{19}$
20	$86^{20} = 4.9 \times 10^{38}$

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- ❖ If we have a figure of *relative* merit, we can classify candidate solutions by how good they are.
- ❖ Heuristic algorithms don't guarantee the ‘best’ solution, but are often adequate (and the only choice!).

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- ❖ **A**nalogy: Trying to find the highest hill by only taking a step uphill from where you are.
- ❖ **P**itfall: Finding a *local* optimum instead of the global optimum.

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- ❖ **Pitfall:** Without good constraints, missing the optimum value.

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- ❖ **Pitfall:** Converging slowly, can still miss best candidate solution. BUT: has a way from getting stuck in local optima.

Example

- ❖ We require:
 - ❑ A problem with relative solution assessment
 - ❑ An algorithm to assess solutions
- ❖ The password cracking didn't have the former.
- ❖ Let's revisit the bag-packing algorithm.

Example

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 - ❑ Brute-force (last lecture)
 - ❑ Hill-climbing
 - ❑ Select heaviest item, then add next heaviest, etc.
 - ❑ Select most valuable item, then add next most valuable item, etc.
 - ❑ Random sampling
 - ❑ Random walk: sample randomly, then iteratively allow change

Setup

```
import numpy as np
import matplotlib.pyplot as plt
import itertools

n = 10
items = list( range( n ) )
weights = np.random.uniform( size=(n,) ) * 50
values = np.random.uniform( size=(n,) ) * 100
```

Setup

```
def f( wts, vals ):  
    total_weight = 0  
    total_value = 0  
  
    for i in range( len( wts ) ):  
        total_weight += wts[ i ]  
        total_value += vals[ i ]  
  
    if total_weight >= 50:  
        return 0  
    else:  
        return total_value
```

Tracking cases

```
max_value = 0.0
max_set = None
lists = []
for i in range(n):
    for set in itertools.combinations( items,i ):
        wts = []
        vals = []
        for item in set:
            wts.append( weights[ item ] )
            vals.append( values[ item ] )
        value = f( wts,vals )
        lists.append( ( wts, value ) )
        if value > 0:
            print( value, wts )
        if value > max_value:
            max_value = value
            max_set = set
```

Tracking cases

```
array = np.array( lists )  
plt.plot( array[:,1], 'b.' )  
plt.xlim( ( 0, len(lists) ) )  
plt.show()
```


Brute-force search

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max_value = 0.0
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        wts = []
        vals = []
        for item in set:
            wts.append( weights[ item ] )
            vals.append( values[ item ] )
        value = f( wts,vals )
        if value > max_value:
            max_value = value
            max_set = set
```

Hill-climbing search

```
max_wt = 50.0
```

```
wts_orig = wts[ : ]
```

```
vals_orig = vals[ : ]
```

```
best_vals = [ ]
```

```
best_wts = [ ]
```

```
best_vals.append( max( vals ) )
```

```
best_wts.append( wts[ vals.index( max( vals ) ) ] )
```

```
wts.remove( wts[ vals.index( max( vals ) ) ] )
```

```
vals.remove( max( vals ) )
```

Hill-climbing search

```
while sum( best_wts ) + wts[ vals.index( max( vals ) ) ] \  
    < max_wt:  
    best_vals.append( max( vals ) )  
    best_wts.append( wts[ vals.index( max( vals ) ) ] )  
    wts.remove( wts[ vals.index( max( vals ) ) ] )  
    vals.remove( max( vals ) )  
  
wts = wts_orig[ : ]  
vals = vals_orig[ : ]
```

Random walk

```
# try a configuration at random
# alter it at random with small likelihood of getting worse
for t in range( 1000 ):
    # two possible moves:  adding or removing
    if f( next_wts,next_vals ) > f( trial_wts,trial_vals ) :
        # if improvement, accept the change
    else:
        # if no improvement, *maybe* accept the change
        # if all-time best, track it
# (see random-walk.py)
```

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- ❖ These run your code many times and return an average time to completion.

Fibonacci sequence

$$F_n = F_{n-1} + F_{n-2} \qquad F_1 = F_2 = 1$$

1 1 2 3 5 8 13 21 34 55 ...

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$$F_n = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^n - \left(\frac{1-\sqrt{5}}{2}\right)^n}{\sqrt{5}}$$

Analytical Fibonacci

```
def fib_a( n ):
    sqrt_5 = 5**0.5;
    p = ( 1 + sqrt_5 ) / 2;
    q = 1 / p;
    return int( (p**n + q**n) / sqrt_5 + 0.5 )
```

Recursive Fibonacci

```
def fib_r( n ):
    if n == 1 or n == 2:
        return 1
    else:
        return fib_r( n-1 ) + fib_r( n-2 )
```

Comparison

```
%timeit fib_a( 12 )  
%timeit fib_r( 12 )
```


Comparison

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%timeit fib_a( 12 )  
%timeit fib_r( 12 )
```

- On my machine, `fib_a` is $55 \times$ faster than `fib_r` for $n = 12$. (Will this performance get better or worse for larger n ?)

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- Which element is compared? It's ambiguous.

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one = np.ones( ( 5, ) )
if one.all() == 1:
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domain = np.linspace( 0,10,11 )
if domain.any() == 1:
    print( 'setup contains one' )
```